

Department

: CSE

Subject Code & Name : CS3492 & Database Management

Semester

PARTICULARS 1 Time Table 2 Student name list 3 Student arrear list 4 Subject Information Record	REMARKS
2 Student name list 3 Student arrear list	
3 Student arrear list	
5 Syllabus	
6 Lesson Plan	
7 Test Plan for the Subject	
8 Result Analysis	
9 Corrective Action Report	<u></u>
10 Quality objective monitoring record	
11 Internal test mark sheet(Consolidated)	
12 Internal test question paper with answer key	
13 Model question paper with answer key	
14 Slip test question paper with answer key	
15 Sample Answer paper for all test(Min-3)	
16 Content beyond the syllabus	
17 Tutorial Class – schedule and content	
18 Assignment – schedule and paper	
19 PPT - handout	
20 Question bank	
21 Sample university question papers(min 5 QP-recent	
22 Personal Log book – Updated	exam)
23 Lecture Note	
zeotare mote	
24 Special Class if any, Approval letter, Schedule, conte	ent covered.

	Prepared By	Approved By
Sign:	0	
Name:		
	Faculty	
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Department : CSE

w.e.f : 08.02.2024

Class	: I	I YEAI	3								Semester :I\	/	
	HOUR	er p	1	11	10.40	111	IV	10.25	v	VI	2.50	VII	VIII
	DAY/ TIME		09.00a.m TO 09.50 a.m.	09.50a.m. TO 10.40a.m.	10.40 a.m TO 10.55 a.m	10.55 a.m. TO 11.45 a.m.	11.45 a.m. TO 12.35p.m.	12.35 p.m. TO 1.20 p.m.	1.20 p.m. TO 2.05p.m.	2.05 p.m. TO 2.50p.m.	2.50 p.m. TO 3.05 p.m.	3.05 p.m. TO 3.50 p.m.	3.50 p.m. TO 4.30 p.m.
MOND	AY	- P. F.				DBMS			2				
TUESD	OAY		DBMS										,
WEDN	ESDAY	11		= ,-	, AK	DBMS		LUNCH			BREAK		
THURS	SDAY				BREAK			רמ	DBMS	DBMS LAB	- B	DBM	IS LAB
FRIDAY	Y		1 ,									DBMS	
SATURI	DAY	-2		·					DBMS				

S.No Sub	ect Code	Name of the Subject	Name of the Staff	hours
	CS3492	DataBase Management System		6
	C55472		TOTAL	6

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Sign:	· *** ********************************		Was
Name:		Mr.S.PRABAKARAN	Dr.M.VIJAYAKUMAR
Time Tab	le Incharge	HOD	Principal



DEPARTMENT OF COMPUTER SCIENCE AND ENGINEERING

DEPARTMENT

: **B.E - CSE**

YEAR / SEM

: II / IV

ACADEMIC YEAR : 2023-2024

S.No	Reg Number	Name of the Student	Dayscholar/ Hostel
1	732422104001	ABISHEK J	Dayscholar
2	732422104002	AKILESH KUMAR S	Dayscholar
3	732422104004	ARUNKUMAR A	Dayscholar
4	732422104005	ASWIN S	Dayscholar
5	732422104006	BASHARATH MAHMOOD S	Dayscholar
6	732422104007	BASKAR S	Dayscholar
7	732422104008	DEEPAK V	Hostel
8	732422104009	DEEPAKRAJ R	Dayscholar
9	732422104010	DHARSHINI R	Dayscholar
10	732422104011	DHARUN T	Dayscholar
11	732422104013	FARGATH A	Hostel
12	732422104014	GUHAN K R	Dayscholar
13	732422104015	GURUPRASAD R	Dayscholar
14	732422104016	HARIJEEVA M	Hostel
15	732422104017	HARIKRISHNAN B	Dayscholar
16	732422104018	HARIPRIYA V	Dayscholar
17	732422104019	IRUDHAYA VISHVA A	
18	732422104020	JEENA D	Dayscholar
19	732422104021	JEEVA S	Dayscholar
20	732422104022	JEEVA S	Dayscholar
21	732422104023	KALAISELVAN R	Dayscholar
22	732422104024	KARTIIIKA K	Dayscholar
23	732422104025	KEERTHIKA S	Hostel
24	732422104027	MAHESWARI T	Dayscholar
25	732422104028	MATHAVAN C	Dayscholar
7,10	732422104030		Dayscholar
26		MOHAMMED THAMEEMMUL ANSARI C J	Dayscholar
27	732422104031	NANDHINI M R	Dayscholar

28	732422104032	NAVANEETHA KRISIINAN M	Dayscholar
29	732422104033	NAVEENA M	Dayscholar
30	732422104034	NAVEEN KUMAR V	Dayscholar
31	732422104035	PANDIE	Dayscholar
32	732422104036	REVATHI P	Dayscholar
33	732422104037	SABARIYANANDIIAN T	Dayscholar
34	732422104038	SARAN B	Dayscholar
35	732422104039	SARAVANAN R	Dayscholar
36	732422104040	SELVAPRIYA C	Dayscholar
37	732422104041	SHANMATHI C T	Dayscholar
38	732422104042	SIKKANTHAR BATHUSHA R	Dayscholar
39	732422104043	SRI RAJ S	Dayscholar
40	732422104044	SUBASH M	Dayscholar
41	732422104045	SWATHI R	Dayscholar
42	732422104046	THIRUPATHI P	Dayscholar
43	732422104047	VASANTH A	Dayscholar
44	732422104048	VASANTHKUMAR P	Dayscholar
45	732422104301	ARUN PRAKASH G	Dayscholar
46	732422104302	DHANUSH B	Dayscholar
47	732422104303	GUNASEKAR S	
48	732422104304	NITHISH KUMAR V	Dayscholar
.,			Dayscholar

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SUBJECT INFORMATION RECORD

Department

CSE

Subject

: Database Management System

Year

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Semester

ÍV

Last year handled by

Percentage of Result (last year)

Quality Objectives

: To produce result more than 80% in University exam

Reference Book

: c.J. Date, A. Kannan,

S. swamynathan, "An

Introduction to Database

Systems", Eight Edition, Pearson Education, 2006.

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Name:		8
	Faculty	HOD

SYLLABUS

CS3492 DATABASE MANAGEMENT SYSTEMS LTPC3003

COURSE OBJECTIVES:

- To learn the fundamentals of data models, relational algebra and SQL
- To represent a database system using ER diagrams and to learn normalization techniques
- To understand the fundamental concepts of transaction, concurrency and recovery
- processing To understand the internal storage structures using different file and indexing techniques
- which will help in physical DB design To have an introductory knowledge about the Distributed databases, NOSQL and database security

UNIT I RELATIONAL DATABASES

10

Purpose of Database System - Views of data - Data Models - Database System Architecture -Introduction to relational databases - Relational Model - Keys - Relational Algebra - SQL fundamentals - Advanced SQL features - Embedded SQL-Dynamic SQL

UNIT II DATABASE DESIGN

8

Entity-Relationship model - E-R Diagrams - Enhanced-ER Model - ER-to-Relational Mapping -Functional Dependencies - Non-loss Decomposition - First, Second, Third Normal Forms, Dependency Preservation - Boyce/Codd Normal Form - Multi-valued Dependencies and Fourth Normal Form - Join Dependencies and Fifth Normal Form

UNIT III TRANSACTIONS

9

Transaction Concepts - ACID Properties - Schedules - Serializability - Transaction support in SQL -Need for Concurrency - Concurrency control -Two Phase Locking- Timestamp - Multiversion -Validation and Snapshot isolation- Multiple Granularity locking - Deadlock Handling - Recovery Concepts - Recovery based on deferred and immediate update - Shadow paging - ARIES Algorithm

UNIT IV IMPLEMENTATION TECHNIQUES

9

RAID - File Organization - Organization of Records in Files - Data dictionary Storage - Column Oriented Storage-Indexing and Hashing-Ordered Indices - B+ tree Index Files - B tree Index Files -Static Hashing - Dynamic Hashing - Query Processing Overview - Algorithms for Selection, Sorting and join operations - Query optimization using Heuristics - Cost Estimation.

UNIT V ADVANCED TOPICS

Distributed Databases: Architecture, Data Storage, Transaction Processing, Query processing and optimization - NOSQL Databases: Introduction - CAP Theorem - Document Based systems - Key value Stores - Column Based Systems - Graph Databases. Database Security: Security issues - Access control based on privileges - Role Based access control - SQL Injection - Statistical Database security - Flow control - Encryption and Public Key infrastructures - Challenges

TEXT BOOKS:

- 1. Abraham Silberschatz, Henry F. Korth, S. Sudharshan, "Database System Concepts", Seventh Edition, McGraw Hill, 2020.
- 2. Ramez Elmasri, Shamkant B. Navathe, "Fundamentals of Database Systems", Seventh Edition, Pearson Education, 2017

REFERENCES:

1. C.J.Date, A.Kannan, S.Swamynathan, "An Introduction to Database Systems", Eighth Edition, Pearson Education, 2006.





(Accredited by NAAC, Under 2f and 12B status)

Faculty Name Department

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: CSE

LESSON PLAN

Designation: Assistant Professor

Semester/ Year: IV/ II

Subject / Code Academic Year : DATABASE MANAGEMENT SYSTEM/CS3492 : 2023-2024

Academ	D.	: 2023-2 oposed	024					
S.No.	Date	Period	Details of Topic Covered	TA	Ref.	Act	tual	Remarks
			UNITT RELATIONAL DATABASES			Date	Period	Kemarks
1	413124	3	Purpose of Database System, Views of data	1	1	413124	2	
·2	5/3/24	1	Data Models	1	1	5/3/24	1)
3	513124	2	Database System Architecture	1	1	513124		
4	63/24	3	Introduction to relational databases	1	1			}
5	6/3/24		Relational Model	1	1	613124		
6	7/3/24	5	Keys, Relational Algebra	1	1	6/3/24	3	
7	7/3124	6	SQL fundamentals	1	1	713124	5	B
8	9/3/24	1	Advanced SQL features	1	1	913124		
9	111324	Ь	Embedded SQL	1	1	11/3/24	'	
10	413/24	-1 -	Dynamic SQL	1	1	14/3/24	Ь	
			UNIT II DATABASE DESIGN			41324	()	* 17
1,5	15/3/24	8	Entity-Relationship model, E-R Diagrams	1	1	15/3/24	8	
2	18/2)24	.3	Enhanced-ER Model, ER-to-Relational Mapping	1		18/3/24		
3	1913/24	1	Functional Dependencies, Non-loss Decomposition	- 1		1913124	3 /	
The second second	22/3/24	7	First normal form, Second normal form	1		22/3/24	1 4	1
5	25/3/24	3.	Third Normal Forms, Dependency Preservation	.1		25/3/24	7	-17L1
6	25/3/24	4	Boyce/Codd Normal Form	1		25/3/24	3	111
7	281324		Multi-valued Dependencies and Fourth Normal Form	1		28/3/24	4 5	8
8	01/4/24	3_	Join Dependencies and Fifth Normal Form	1				. 3.
#2			UNIT III TRANSACTIONS			114/24	3 /	
-1	01/4/24	5	Transaction Concepts , ACID Properties	1	1 6	1	= 0	
-	05/4/24	7	Schedules , Serializability	1	-		5()	
	08/4/24	3	Transaction support in SQL, Need for Concurrency	1	-		7 4	
	09/4/24		Concurrency control ,Two Phase Locking-	1		814124 3		-
5	13/4/24	1	Timestamp , Multiversion			1424	10	7
6	13/4/24	5	Validation and Snapshot isolation—,Multiple Granularity locking			3/4/24 5	18	1
7	514124	3	Deadlock Handling ,Recovery Concepts	-		14/24 3	1	

8	264124	7	Recovery based on deferred and immediate update	1	1	26/4/24	7	
9	27(4)24	6	Shadow paging , ARIES Algorithm	1	1	27/4/24	Ь	
			UNIT IVIMPLEMENTATION TECHNIQUES	5				
1	29/4/24	3	RAID - File Organizatio	1	1	29/4/24	3	
2	3014124	1	Organization of Records in Files ,Data dictionary Storage	1	1	3014/24	1	
3	06/5/24	3	Column Oriented Storage, Indexing and Hashing	1	1	06/5/24	3	
4	01/5/24	1 .	Ordered Indices , B+ tree Index Files	1	1	07/5/24	1 -	130
5	0915124	5	B tree Index Files, Static Hashing	1	1	09/5/24	5	
6	09/5/24	6	Dynamic Hashing ,Query Processing Overview	1	1	09/5/24	6	4
7	10/5/24	7	Algorithms for Selection	1	1	1015124	+	
8	11/5/24	5	Sorting and join operations	1	1	1115124	5	*
9	13/5/24	3	Query optimization using Heuristics, Cost Estimation	1	1	13/5/24	3	
			UNIT V ADVANCED TOPICS				Q.	/
1	17/5/24	7	Distributed Databases: Architecture, Data Storage	1	1	1715/24	7	
2	2015/24	3	Transaction Processing, Query processing and optimization	1	1	2015/24	3	1 Y
3	21/5/24	1	NOSQL Databases, Introduction , CAP Theorem , Document Based systems	1	1	21/5/24	1-7-	-
4	22/5/24	3	Key value Stores, Column Based Systems	1	1	22/5/24	3	
5	23/5/24	5	Graph Databases	1	1	23/5/20		T.
6	27/5/24	3	Database Security, Security issues	1		27/5/24		
			Access control based on privileges, Role Based access					
7	2815124	1	control	1	1	2815/24	1,.	
7	2815124 2915124	3	. •	1	1	28/5/24 29/5/24	-	40

Reference books (Ref):

1.. C.J.Date, A.Kannan, S.Swamynathan, "An Introduction to Database Systems", Eighth Edition, Pearson Education, 2006.

Teaching Aids (TA):

- 1. Black Board with Chalk
- 2. Overhead Projector
- 3. LCD Projector
- 4. Others (Field vists, Charts, Cutset Models)

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Name: Faculty	ROD	Principal



TEST PLAN FOR SUBJECT

Subject Database Management System Faculty: Stratakoran

Semester

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Year:

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Department:

CSE

S. No.	Description	Planned Date/Month	Actual Conducted Date / Month	Remarks
Ι.	Unit test-I	27/3/2024	27/3/2024	
2,	Unit test-ji	17/4/2024	Huloon	
3.	Unit test- 111	11/5/2024		
t.	Unut test - i	941 21 2024	24/5/2024	
	Model Exam-7	01/06/2024	01/06/2024	有差 力

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Subject

: Database Management System Date: 2.4. 2024

Class

Year

Department : CSE

Semester

Exam details & date

: Unit test-1 & 27/3/24

Faculty

: S. Prabakaran

Number of students

7+3

No. of students attended

33

No. of students absent

14

No. of students passed

33

No. of students failed

Percentage of failures

01.

Marks	0-25	26-50	51-75	76-90	91-100
No. of Students	,-	33			

	Prepared By	Approved By
Sign:	Q /	
Name:		E
	Faculty	HoD



Subject

: Database Management System Date: 19.4.24

Class

: II Year

Department : CSE

Semester

ĺV

Exam details & date

: Unit test - 1 & 17/4/24

Faculty

: S. Prabacanan

Number of students

: 47

No. of students attended

: 46

No. of students absent

1

No. of students passed

: 38

No. of students failed

8

Percentage of failures

: 17.1.

Marks	0-25	26-50	51-75	76-90	91-100
No. of Students	8	38		- 5	3

	Prepared By	Approved By
Sign:		0 -
Name:		
	Faculty	
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Subject

: Database Management System Date: 14/5/24

Class

li year

Department : CSE

Semester

Exam details & date

: Unit test - 111 & 1115124

Faculty

: S. Prabakaran

Number of students

47

No. of students attended

45

No. of students absent

2

No. of students passed

42

No. of students failed

Percentage of failures

6.61.

Marks	0-25	26-50	51-75	76-90	91-100
No. of Students	3	42			

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Sign:	8	
Name:		
	Faculty	HoD



Subject

: Database Management System Date: 27/5/24

Class

il year

Department : CS

Semester

IV

Exam details & date

Unit test - 1 & 24/5/24

Faculty

: S. Prabakaian

Number of students

: 47

No. of students attended

: 47

No. of students absent

No. of students passed

: 46

No. of students failed

: 1

Percentage of failures

: 2.1.1.

					a d
Marks	0-25	26-50	51-75	76-90	91-100
No. of Students)	46		29	

	Prepared By	A
Sign:		Approved By
Name:		
	Faculty	
A - 10		HoD



Subject

: Database Management System Date: 4/6/24

Class

Year

Department : CSE

Semester

Exam details & date

: Model Exam - ? & oilob124

Faculty

: S. prabakaran

Number of students

47

No. of students attended

: 47

No. of students absent

No. of students passed

: 43

No. of students failed

Percentage of failures

: 8.5%

Marks	0-25	26-50	51-75	76-90	91-100
No. of Students	-	4	25	16	2

	Prepared By	Approved By
Sign:	War and the second seco	0
Name:		
	Faculty	HoD



CORRECTIVE ACTION REPORT

Dept.

: CSE

Year

Subject

: Database Management Semester

ī	1	D	()	stem		
S.No	Unit Test	Percentage of marks	Root Cause (Metrics)	Corrective Action	Deadline date	Remarks
	Unit test I	100 1.	•	-	_	
2	Unil E Lest ij	831.	_	_	_	
3	Unit test in	931.	-	_	_	-
4	Unil testily	97.	<u> </u>	_		
5	Model -I	917.	-	<u>-</u>		_

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Name:			
	Faculty	HOD	



QUALITY OBJECTIVE MONITORING RECORD

Department : CSE

Year

Semester

Subject

: Database Management System

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S.N 0	Quality Objective	Expecting result	Obtained result	Expecting result	Obtained result								
	To Produce result					7							
	more than \$0+. in University Exam	30 У.	loot	90%	834	qa).	934	da.	974	907	914		
TOTAL TO BE TO A SERVICE AND A	Exam										Acceptance of the control of the con		There is no extra control of the con

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Name:		
	Faculty	HOD



DEPARTMENT : B.E - CSE

YEAR / SEM : II / IV

ACADEMIC YEAR : 2023-2024

STUDENT MARK LIST

		STUDENT MARK	LIST		I		1
S.N o	Reg Number	Name of the Student	Unit Test -	Unit Test - 2	Unit Test -	Unit Test - 4	Mode test - 1
1	732422104001	ABISHEK J	AB	36	AB	36	52
2	732422104002	AKILESH KUMAR S	AB	43	39	39	83
	732422104004	ARUNKUMAR A	12	31	35	34	30
3	732422104005	ASWINS	26	37	32	35	73
4	732422104006	BASHARATH MAHMOOD S	34	39	36	36	72
5	732422104007	BASKAR S	31	45	44	33	78
6	732422104008	DEEPAK V	30	32	AB	31	12
7	732422104009	DEEPAKRAJ R	32	29	22	34	71
8		DHARSHINI R	47	47	45	39	91
9	732422104010	DHARUN T	33	42	35	31	62
10	732422104011	FARGATH A	40	44	35	36	74
11	732422104013	GUHAN K R	29	34	31	39	61
12	732422104014	GURUPRASAD R	AB	35	34	37	77
13	732422104015	HARIJEEVA M	35	39	34	37	56
14	732422104016	HARIKRISHNAN B	37	42	35	39	77
15	732422104017	HARIPRIYA V	46	41	38	45	77
16	732422104018	IRUDHAYA VISHVA A	39	41	42	39	84
17	732422104019		39	35	39	37	82
18	732422104020	JEENA D	AB	AB	29	31	47
19	732422104021	JEEVA S	AB	39	29	34	59
20	732422104022	JEEVA S	31	42	29	36	78
21	732422104023	KALAISELVAN R	44	41	37	37	82
22	732422104024	KARTHIKA K	34	35	27	21	50
23	732422104027	MAHESWARI T	39	36	41	34	80
24	732422104028	MATHAVAN C MOHAMMED THAMEEMMUL	AB	34	24	36	81
25	732422104030	ANSARI C J			28	34	74
26	732422104031	NANDHINI M R	40	36	23	30	69
27	732422104032	NAVANEETIIA KRISIINAN M	AB	29	31	38	81
28	732422104033	NAVEENA M	36	34		27	71
29	732422104034	NAVEEN KUMAR V	34	24	24	36	70
30	732422104035	PANDI E	31	23	4	37	88
31	732422104036	REVATHI P	45	46	44	29	72
32	732422104037	SABARIYANANDHAN T	35	32	33	37	81
33	732422104038	SARAN B	39	39	39		75
	732422104939	SARAVANAN R	31	37	31	32	91
34	732422104040	SELVAPRIYA C	43	45	45	42	-/1

36	732422104041	SHANMATHI C T		*			
37	732422104042	Marine distribution of the Control o	AB	39	39	38	79
38	732422104043	SIKKANTHAR BATHUSHA R	32	35	34	26	72
39		SRI RAJ S	36	27	42	35	-
	732422104044	SUBASH M	AB	29	31		75
40	732422104045	SWATHI R	34	5		29	54
41	732422104046	THIRUPATHIP			34	26	47
42	732422104047	VASANTILA	AB	21	36	27	43
43	732422104048	VASANTHKUMAR P	AB	, 29	34	25	49
44	732422104301	ARUN PRAKASH G	34	21	29	27	78
15	732422104302		AB	15	27	30	71
-		DHANUSH B	28	22	29	32	59
16	732422104303	GUNASEKAR S	AB	24	31	26	
17	732422104304	NITHISH KUMAR V	AB	27	37	35	62 71

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NAME			
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ggrod to	n Xa ² -	UNIT TEST	I Later	Date/Session	27.03.2024	Marks		50
Course cod	e	CS3492	Course Title	Database Mana	gement System			
Regulation		2021	Duration	1:30 Hours	Academic	Year	2023-2024	
Year		II	Semester	IV	Departme	nt	CSE	
COURSE	OUTCO	MES						
CO1:	Cons	truct SQL Querie	es using relational a	algebra				
CO2:	Desig	gn database using	ER model and not	rmalize the database	;			
CO3:	Cons	truct queries to h	andle transaction p	rocessing and maint	tain consistency	y of the da	atabase	
CO4:	Com	pare and contras	t various indexing	strategies and apply	the knowledge	e to tune th	he perform	mance
	of the	database						
CO5:	Appr	aise how advance	ed databases differ	from Relational Da	tabases and fin	d a suitab	le databa	se for
	the g	iven requirement	•					

0.17	Question	CO	BTS
Q.No.	PART A		
	(Answer all the Questions 10 x 2 = 20 Marks)	CO1	R
1	Define data abstraction and different levels of data abstraction.	CO1	K
		C01	R
2	What is DBMS ?Why do we need a DBMS?		1
3	Differentiate between Dynamic SQL and static SQL?	CO1	A
		CO1	U
4	What is different primary key and foreign key?		
5	Name the categories of SQL commands.	CO1	R
5		CO1	U
6	What is the data definition language? give example		
	List the reason why null value might be introduced into the database.	C01	R
7	List the reason why null value inight be introduced.	CO1	A
8	Why key is essential? & write the different types of keys.		R
<u>9</u>	What is meant by instance and schema of the database?	C01	
	what is meant by instance and conventional file system.	CO1	A
10	Different between database system and conventional file system.	Annual Page 15	
	PART B		
	(Answer all the Questions 2 x 15 = 30 Marks)	COI	F
11a	Write about database system Architecture with neat diagram.		
11-	OK	CO1	
11b	Briefly explain about Database system architecture	CO1	
	Explain select project, Cartesian product and join operations in relational	CO1	A
12a_	OR	A Partie	-
		COI	(
12b	Explain the codd's rules for relational design.	Jan Carlo	

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A CONTRACTOR	UNIT TE	STI	Date/Session	27.03.2024	Mar	ks	50
Course code	CS3492	Course Title	Database Man	agement Syster	n		
Regulation	2021	Duration	1:30 Hours	Academic	Year	2023-202	24
Year	111	Semester	IV	Departme	ent	CSE	

PART A

1. Define data abstraction and different levels of data abstraction.

- Data abstraction is a technique in database systems to simplify complex data by hiding lowlevel details and focusing on essential aspects of data for easier management and interaction.
- There are three levels of data abstraction:
 - Physical level: The lowest level, describing how data is stored on the physical storage media.
 - o Logical level: Describes the structure of the database and relationships, without detailing how data is stored.
 - o View level: The highest level, defining different views or ways data can be seen by different users based on access needs.

2. What is DBMS?

DBMS stands for **Database Management System**. It is a software that enables the creation, management, and use of databases, providing an interface for users to interact with stored data.

It helps in organizing data, ensuring data integrity, security, and supporting transactions, making it easier for users to store, retrieve, and manage data efficiently.

Why do we need a DBMS?

- A Database Management System (DBMS) is essential for managing large volumes of data in an organized and efficient manner, allowing for storage, retrieval, and manipulation of data.
- DBMS ensures data security, integrity, and ease of access, along with providing functionalities like backup, recovery, and concurrency control.

3. Differentiate between Dynamic SQL and Static SQL.

- Dynamic SQL: SQL statements that are constructed and executed at runtime, allowing for more flexibility and handling of various scenarios, as the SQL code can adapt to changing inputs.
- Static SQL: SQL statements written and embedded in the application code, known at compile-time, offering better performance and security as the SQL commands are fixed.

4. What is the difference between primary key and foreign key?

- Primary Key: A unique identifier for each record in a table, ensuring that no two rows have
 the same value for this field.
- Foreign Key: A field in one table that is a primary key in another table, used to establish a
 relationship between the two tables.

5. Name the categories of SQL commands.

- Data Definition Language (DDL): Commands that define and structure database objects, such as CREATE, ALTER, and DROP.
- Data Manipulation Language (DML): Com nands that handle data operations, such as INSERT, UPDATE, and DELETE.
- Data Control Language (DCL): Commands that manage access permissions, such as GRANT and REVOKE.
- Transaction Control Language (TCL): Commands that manage transaction operations, such as COMMIT and ROLLBACK.

6. What is the data definition language? Give an example.

- Data Definition Language (DDL) is a set of SQL commands used to define and modify the structure of database objects, like tables and indexes.
- Example: CREATE TABLE students (id INT PRIMARY KEY, name VARCHAR(50));

7. List the reasons why a null value might be introduced into the database.

- The value is unknown or missing for a certain attribute in a record.
- The attribute does not apply to the specific record or entity, hence leaving it empty (e.g., "middle name" for someone without one).

8. Why is a key essential in a database?

- Keys are essential for uniquely identifying each record in a table, ensuring data integrity and enabling efficient retrieval of records.
- They also establish relationships between tables, allowing for the linking of related data across multiple tables.

What are the different types of keys?

- Primary Key: Unique identifier for each row in a table.
- Foreign Key: Establishes a relationship between two tables.
- Candidate Key: Any attribute or set of attributes that can uniquely identify a row.
- Super Key: A combination of attributes that uniquely identifies a row.
- Composite Key: A key that consists of more than one attribute to uniquely identify a row.
- Alternate Key: A candidate key that was not chosen as the primary key.

9. What is meant by instance and schema of the database?

- Instance: Refers to the actual data stored in a database at a specific point in time; it is
- Schema: The overall structure or design of the database, defining tables, fields, and relationships; it is relatively static and only changes if the database structure is altered.

10. Differentiate between database systems and conventional file systems.

- Database Systems: Organized collections of data managed by DBMS, offering features like
 Conventional File Systems: Organized collections of data managed by DBMS, offering features like
- Conventional File Systems: Simple storage systems without DBMS features, lacking centralized data management, often resulting in data redundancy and inconsistency

11a. Write about Database System Architecture with a Neat Diagram.

Database system architecture defines how different components of a Database Management System (DBMS) interact to manage, store, process, and retrieve data. It plays a crucial role in designing efficient database systems by organizing the interaction between users, applications, and data storage in a structured way.

The three-tier architecture is a commonly used structure that divides the database system into three main layers:

Presentation Layer (User Interface Layer):

- 1. This is the topmost layer, where end-users interact with the system through a graphical user interface (GUI) or an application.
- 2. This layer provides interfaces that allow users to perform operations such as querying, updating, and retrieving data.
- 3. Typically, the presentation layer doesn't interact directly with the database; instead, it forwards requests to the application layer.
- 4. By keeping the user interface separate, this layer allows easy modifications in the appearance and user interactions without impacting the database.

2.

Application Layer (Business Logic Layer):

- 1. The application layer, or middle layer, serves as the intermediary between the user interface and the database.
- 2. It contains the core logic for processing user requests and ensuring data accuracy, often implementing validation, data processing, and error handling.
- 3. This layer transforms user inputs into database operations and enforces business rules. For example, it may check if an item is in stock before processing a purchase order.
- 4. This separation allows enhanced security since direct access to data is restricted, and only authorized requests can modify or access the database.
- 5. Additionally, it improves scalability, as the application layer can be adjusted to handle growing numbers of users or increased data processing needs.

3.

Database Layer (Data Storage Layer):

- 1. The database layer is where the data resides, organized and managed by the DBMS.
- 2. This layer is responsible for storing data physically on the storage media and managing access through indexing, query processing, and transaction handling.
- 3. Features of the DBMS, such as data consistency, concurrency control, and recovery mechanisms, operate here to maintain data integrity.
- 4. The database layer can also provide optimizations like indexing and caching to improve query performance, particularly in large-scale systems.
- 5. By keeping the data management isolated, this layer allows the application and presentation layers to function independently of the physical data storage details

Diagram:

Presentation Layer (User Interface) | Application Layer (Business Logic) Database Layer (Data Storage)

11b. Briefly Explain Database System Architecture

Database system architecture organizes the interaction between different components in a DBMS, creating a structured approach for data management. This structure ensures efficient data processing, secure access, and system reliability by separating user interfaces, application logic, and data storage.

Types of Database System Architectures:

Two-tier Architecture:

- 1. In this simpler setup, the client application communicates directly with the database server.
- 2. Client Layer: Users interact with the application on this layer, sending requests and receiving data directly from the database.
- 3. Database Layer: The DBMS manages data storage, retrieval, and integrity, directly responding to client requests.
- 4. This type is suitable for small-scale applications where direct access provides faster performance with limited users.
- 5. However, it lacks the ability to scale well for large user bases and may present security risks since each client has direct access to the database.

Three-tier Architecture:

- 1. Three-tier architecture introduces an application layer between the user interface and database, making it a preferred choice for larger applications.
- 2. Presentation Layer (User Layer): The top layer where users interact with the system, often through web browsers, mobile apps, or GUIs.
- 3. Application Layer (Business Logic Layer): Acts as an intermediary that processes user requests, enforces business rules, and validates inputs before passing them to the database. This layer also handles responses back to the user, adding security by managing data access.
- 4. Database Layer (Data Layer): The lowest layer responsible for storing, retrieving, and organizing data. The DBMS in this layer manages transactions, indexing, and backup, among other functions.
- 5. This architecture is more secure and scalable. Since the database is isolated, it's protected from unauthorized access, and the middle layer allows for consistent data handling and validation.

Benefits of Three-Tier Architecture:

- Enhanced Security: Direct access to the database is restricted by the application layer,
- Scalability: Additional users and resources can be managed more effectively since the application layer can be scaled independently.
- Maintainability: Changes in business rules or database structure can be managed in the application layer without impacting other components.
- Flexibility: Modifications in the user interface or data structure do not affect each other directly, allowing for better data independence.

In relational algebra, these fundamental operations help in retrieving and manipulating data from relational databases by performing various transformations on tables (relations). Each operation has a specific role in filtering, combining, or restructuring data.

Operations in Relational Algebra:

Select (v):

- 1. The Select operation is used to filter rows from a relation based on a given condition. Only rows that satisfy the condition are included in the result.
- 2. Notation: σ condition (Relation)
- 3. Example: If we have an Employees relation and want to retrieve employees older than 25, we write σ age > 25 (Employees). This will return only rows where the age attribute is greater than 25.
- 4. Usage: Select helps narrow down data based on specific criteria, making it easier to work with relevant subsets of a larger dataset.

Project (π) :

- 1. The Project operation is used to select specific columns (attributes) from a relation, essentially reducing the number of columns in the output.
- 2. Notation: π attributes (Relation)
- 3. Example: π name, age (Employees) retrieves only the name and age columns from the Employees table, omitting other columns.
- 4. Usage: Project is useful when only certain attributes are needed for analysis or reporting, reducing the volume of data retrieved.

Cartesian Product

- 1. The Cartesian Product operation combines two relations by pairing every row from the first relation with every row from the second. It is a foundational operation for other complex operations, like join.
- 2. Notation: Relation1 × Relation2
- 3. Example: If we have Employees and Departments, Employees × Departments results in a new relation with all possible combinations of employees and departments.
- 4. Usage: The Cartesian Product is used primarily as an intermediate step for joins but can be computationally intensive due to the large number of combinations generated.

Join

- 1. The Join operation combines rows from two relations based on a related attribute or condition, matching rows with common values.
- 2. Types of Joins:
 - 1. Theta Join: Matches rows based on a condition (☐ condition).
 - 2. Equi-Join: A join with equality condition (=).
 - 3. Natural Join: Matches rows with the same value in common attributes, eliminating duplicate columns.
- 3. Notation: Relation1 condition Relation2
- 4. Example: Employees Employees.department_id = Departments.department_id Departments combines employee and department data where department_id matches.
- 5. Usage: Joins are essential for combining related data from multiple tables, allowing complex queries across related entities in a database.

12b. Explain Codd's Rules for Relational Design.

E.F. Codd, the pioneer of the relational database model, formulated 12 rules (commonly called "Codd's Rules") to define what constitutes a relational database and how it should be designed. These rules ensure that a database fully utilizes relational principles for efficient, consistent, and flexible data management.

Codd's 12 Rules for Relational Database Systems:

Rule 0 (Foundation Rule):

- 1. A relational database management system (RDBMS) must manage data entirely through its relational capabilities, adhering to all other rules.
- 2. This rule emphasizes that an RDBMS must fundamentally be relational.

Rule 1 (Information Rule):

- 1. All information in a database should be stored in tables (relations), with rows and columns.
- 2. This structure allows for a consistent, organized way of storing data.

Rule 2 (Guaranteed Access):

- 1. Each value in a table must be accessible by specifying the table name, primary key, and column name.
- 2. This rule ensures that data can be retrieved using a straightforward method, regardless of table complexity.

Rule 3 (Systematic Treatment of Nulls):

- 1. Null values must be supported to represent missing or inapplicable data, allowing for three states: true, false, and unknown.
- 2. This flexibility accommodates situations where data is incomplete or unknown.

Rule 4 (Dynamic Online Catalog):

- The database's metadata (catalog) should also be stored in tables and be accessible via SQL.
- 2. Users can query database structure information in the same way as data, supporting transparency.

Rule 5 (Comprehensive Data Sub-language):

- 1. The RDBMS should support a powerful data sub-language, such as SQL, for defining, manipulating, and controlling data and transactions.
- 2. This ensures a complete command set for interacting with the database.

Rule 6 (View Updating Rule):

- Users should be able to modify data through views, as long as such modifications do not affect the underlying integrity.
- 2. This feature allows flexibility in data presentation while preserving data accuracy.

Rule 7 (High-Level Insert, Update, and Delete):

- 1. The system must support set-based operations, allowing rows to be inserted, updated, or deleted in bulk rather than one at a time.
- 2. This capability improves efficiency in data manipulation, especially for large datasets.

Rule 8 (Physical Data Independence):

1. The application should remain unaffected by changes in how data is physically stored.

2. This rule supports ease of optimization and hardware changes without affecting applications.

Rule 9 (Logical Data Independence):

1. The application should remain unaffected by changes to the logical structure of the database.

2.Logical data independence protects applications from schema modifications like adding new fields.

Rule 10 (Integrity Independence):

1. This ensures that data rules are defined at the database level, maintaining consistency and security.

2.Logical data independence protects application from schema modifications like adding new fields

Rule 11 (Distribution Independence):

1. The RDBMS should function correctly regardless of data being distributed across multiple locations.

2. This rule supports flexibility for distributed databases, enabling data to be accessed or modifieds scamlessly

Rule 12 (Non-Subversion Rule):

1. Low-level access methods should not bypass or undermine the RDBMS's integrity constraints or rules.

2. This rule prevents unauthorized modifications, ensuring data integrity and consistency.

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71	UNIT TI	EST II	Date/Session	17.04.2024	Marks	50	
Course co	de CS3492	Course Title	Database Mana	gement System			
Regulation 2021 Duration		1:30 Hours	Academic 1	demic Year 2023-2			
Year	II	Semester	IV	Department		CSE	
COURSE	OUTCOMES						
CO1:	Construct SQL Q	ueries using relational a	lgebra				
CO2:	Design database u	sing ER model and non	nalize the database				
CO3:		to handle transaction pr			of the dat	tabase	
CO4:		trast various indexing st					
	of the database			•		•	
CO5:	Appraise how adv	anced databases differ f	rom Relational Data	bases and find	a suitable	database for	
,	the given requiren						

Q.No.			BT	
	PART A (Answer all the Questions 10 x 2 = 20 Marks)			
1	Explain Entity relationship model.	CO2	R	
2	Give the limitation of E-R model? How do you overcome this?	CO2	A	
3	What is an entity?	CO2	R	
4	Define functional dependency?	CO2	R	
5	Why certain functional dependencies are called trival functional dependency?	CO2	R	
6	Define normalization.	CO2	U	
7	What is multivalued dependency?.	CO2	R	
	Describe BCNF and describe a relation which is in BCNF?	CO2	A	
8	Why it is necessary to decompose a relation?	CO2	A	
10	Explain at least two desirable properties of decomposition.	CO2	R	
May I	PART B			
1	(Answer all the Questions 2 x 15 = 30 Marks)	CO2	E	
11a	Explain in detail about normal forms. OR			
	CED Hoursm with example Briefly.	CO2	E	
11b	Explain the various components of ER diagram with example Briefly.	CO2	C	
12a	Explain the various components of ER diagrams Explain the various components of ER diagrams Explain the various components of ER diagrams OR			
	- OK	CO2	U	
12b	Discuss about integrity constraints in SQL?			

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	UNIT TE	ST II	Date/Session	17.04.2024	Marks	50	
Course code	CS3492	Course Title	Database Management System				
Regulation	2021	Duration	1:30 Hours	Academic Yea	-2024		
Year II		Semester	IV	Department	CSE	CSE	

PART A

1. Explain Entity relationship model.

The Entity-Relationship Model (ER Model) is a conceptual framework used in database design to visually represent the structure and relationships of data within a system.

- 2. Give the limitation of E-R model? How do you overcome this?
 - Limited Representation of Data Constraints
 - Difficulty Handling Complex Relationships
 - Lack of Support for Temporal Data
 - Limited Scalability for Complex Systems
 - Inflexibility in Representing Advanced Data Types

3. What is an entity?

An entity in the context of database design and the Entity-Relationship (ER) model is a distinct object, concept, or "thing" that can be identified and represented within a system. Entities represent real-world items or abstractions that need to be stored and managed in a database.

4.Define functional dependency.

Functional dependency is a relationship in which one attribute uniquely determines another attribute within a relation. For example, if attribute A determines attribute B, knowing A's value means we can uniquely identify B's value. Functional dependencies help establish rules for data integrity and are critical in database normalization.

5. Why are certain functional dependencies called trivial functional dependencies?

Trivial functional dependencies occur when an attribute functionally determines itself, such as $\{A\} \rightarrow A$, or when one subset determines another subset within the same attributes. These dependencies are inherently satisfied, as the determinant set contains or equals the determined attribute. Trivial dependencies are always present but do not affect normalization.

6.Define normalization.

Normalization is the process of organizing data in a database to reduce redundancy and improve data integrity. It involves dividing tables into smaller relations and establishing dependencies to minimize anomalies during data manipulation. The goal of normalization is to create a structure that supports efficient storage and accurate retrieval of data.

7. What is a multivalued dependency?

A multivalued dependency exists when an attribute in a relation uniquely determines multiple independent values of another attribute, creating a one-to-many relationship. This situation leads to redundancy

if not managed properly. Multivalued dependencies are common in scenarios with repeating data, such as lists of skills for a person.

8.Describe BCNF and describe a relation which is in BCNF.

Boyce-Codd Normal Form (BCNF) is an advanced version of the Third Normal Form. A relation is in BCNF if every determinant is a candidate key, ensuring no anomalies related to functional dependencies. For example, a relation where each attribute depends only on a candidate key is considered in BCNF, eliminating redundancy more effectively than 3NF.

9. Why is it necessary to decompose a relation?

Decomposition is necessary to eliminate redundancy and prevent data anomalies (insert, update, delete) in a database. By breaking down a relation, it ensures data integrity, maintains a logical structure, and simplifies queries. Proper decomposition helps optimize storage and minimizes data inconsistency in relational databases.

10.Explain at least two desirable properties of decomposition.

Two desirable properties of decomposition are Lossless Join and Dependency Preservation. Lossless join ensures that decomposed tables can be joined back together without losing information, maintaining data completeness. Dependency preservation means that functional dependencies are preserved in the new tables, making constraint enforcement simpler.

PART B

11a. Explain in detail about Normal Forms.

First Normal Form (1NF):

- o Ensures that each attribute contains atomic (indivisible) values.
- o Removes repeating groups, making sure each column contains unique values.
- o Example: If a table has multiple phone numbers in one cell, split each into separate rows.

Second Normal Form (2NF):

- o Builds on 1NF and ensures that all non-key attributes depend on the entire primary key.
- o Applies to tables with composite keys, eliminating partial dependencies.
- Example: In an "Orders" table with {OrderID, ProductID}, remove attributes like "CustomerName" that depend on only one part of the key.

Third Normal Form (3NF):

- o Ensures no transitive dependencies, where non-key attributes depend on other non-key attributes.
- o All attributes must depend directly on the primary key.
- Example: In a "Students" table, move "Department" details to a separate "Department" table if they depend on "DepartmentID."

Boyce-Codd Normal Form (BCNF):

- o A stricter form of 3NF where every determinant is a candidate key.
- o Resolves cases where 3NF is insufficient due to dependencies that violate candidate keys.
- Example: If a professor teaches only one subject, split the "Professor" and "Subject" tables if dependencies are unclear.

Fourth Normal Form (4NF):

- Removes multivalued dependencies, ensuring that attributes are independent.
- Deals with situations where one attribute can determine multiple values of another attribute independently.
- Example: If "StudentID" determines both "Course" and "Hobby" independently, separate them into distinct tables.

11b. Explain the various components of ER Diagram with examples briefly.

Entities:

- Objects or concepts that have a distinct existence in the database.
- o Represented by rectangles, such as "Employee" or "Department."

Attributes:

- o Characteristics or properties of entities.
- Shown as ovals, for example, "EmployeeName" for an "Employee" entity.

Relationships:

- Connections between two or more entities.
- o Represented by diamonds, such as a "WorksFor" relationship between "Employee" and

Primary Key:

- o A unique attribute identifying each entity instance.
- o Example: "EmployeeID" is the primary key for "Employee."

Cardinality:

- Specifies the number of instances in one entity that can relate to instances in another.
- o Example: 1-to-many between "Department" and "Employee."

12a. Explain in detail about Multivalued Dependencies and Fourth Normal Form.

Multivalued Dependency (MVD):

- o Occurs when one attribute uniquely determines multiple independent values of another attribute. o Causes redundancy if not resolved, as one attribute controls multiple values.
- o Example: In a "Student" table with "StudentID" determining both "Course" and "Hobby" independently, "StudentID → Course" and "StudentID → Hobby" are MVDs.

Fourth Normal Form (4NF):

- Builds on BCNF and eliminates multivalued dependencies.
- o Ensures no attribute has multiple independent values dependent on the same primary key.
- Example: Split "Student" into two tables, one for "StudentID" and "Course" and another for

2b. Discuss about integrity constraints in SQL?

Integrity constraints in SQL ensure the accuracy and consistency of data within a relational database. These constraints define rules for permissible values in tables to maintain the correctness and reliability of the lata. Here's an overview of common integrity constraints in SQL:

L Primary Key Constraint

A Primary Key uniquely identifies each row in a table. It must contain unique values, and it cannot contain NULL values.

- Purpose: Ensure each record is unique and identifiable.
- Example

```
CREATE TABLE Employee (
EmplD INT PRIMARY KEY,
Name VARCHAR(50),
DeptID INT
);
Here, EmplD is the primary key, meaning each EmplD in the Employee table must be unique and not null.
```

Foreign Key Constraint

A Foreign Key establishes a relationship between two tables by linking a column in one table to the primary key in another table.

• Purpose: Maintain referential integrity between tables.

```
Example
CREATE TABLE Department (
DeptID INT PRIMARY KEY,
DeptName VARCHAR(50)
);

CREATE TABLE Employee (
EmpID INT PRIMARY KEY,
Name VARCHAR(50),
DeptID INT,
FOREIGN KEY (DeptID) REFERENCES Department(DeptID)
);
In this example DeptID in the Employee table is a foreign key that
```

In this example, DeptID in the Employee table is a foreign key that references the primary key DeptID in the Department table, ensuring that each employee belongs to a valid department.

Unique Constraint

The Unique constraint ensures that all values in a column or a set of columns are distinct across the table. Unlike the primary key, a table can have multiple unique constraints, and unique columns can contain a single NULL value (in most SQL implementations).

Purpose: Prevent duplicate values in specific columns.

Example

```
CREATE TABLE Employee (
  EmpID INT PRIMARY KEY,
 Name VARCHAR(50) UNIQUE.
  DeptID INT
):
```

Not Null Constraint

The Not Null constraint ensures that a column cannot contain NULL values. It is often used to ensure required fields are filled in.

- Purpose: Ensure essential data is not missing.
- Example

CREATE TABLE Employee (EmpID INT PRIMARY KEY, Name VARCHAR(50) NOT NULL, DeptID INT NOT NULL);

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Comme		UNIT TEST	III	Date/Session	11.05.2024	Mark				
		CS3492	Course Title				(5	50		
Regulation		2021	Duration	Database Mana						
Year		II		1:30 Hours	Academie	Year	2023-20	024		
COURSE OUTC			Semester	IV	Department		CSE			
CO1:			rios usina -1 1					4 7.16		
CO2:	Construct SQL Queries using relational algebra Design database using ER model and normalize the database									
CO3:	Con	struct succional	g ER model and nor	malize the database						
CO4:	Con	struct queries to	handle transaction p	rocessing and maint	ain consistency	of the d	atabase			
	of th	npare and contra le database	st various indexing	strategies and apply	the knowledge	to tune t	the perfe	ormance		
CO5:	App		ced databases differ	from Relational Dat						

Q.No.	Question	СО	DT
	PART A (Answer all the Questions 10 x 2 = 20 Marks)	1 0	BT
1	What is transaction.	CO3	R
2	List out ACID property?		
3		CO3	R
	Why is it necessary to have control of concurrent execution of transaction? how it is possible?	CO3	R
4	What are the different modes of lock?	CO3	U
5	Define deadlock.	СО	R
6	List the four conditions for dead lock		
7		CO3	U
	What is meant by concurrency control.	CO3	R
8	Define two phase locking protocol.	CO3	-
9	What is meant by serialization? How it is tested?	CO3	A
10		CO3	R
10	Different strict two phase locking and rigourous two phase locking protocol.	CO3	R
	PART B		y Carry
11a	(Answer all the Questions 2 x 15 = 30 Marks) Discuss the violation caused by each of the C H		
	Discuss the violation caused by each of the following: lost update problem, dirty read, non repeatable read and phantoms with suitable example.	CO3	E
	OP		1
11b	Explain Deadlock in detail with an example.	CO3	E
12a	(1) Write short notes on transition, concepts and ACID	CO3	A
	(ii) Explain conflict serializability and view serializability.	COS	
12b	OB		76
120	Define functional dependencies. How are primary keys related to FD's?	CO3	C

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UNIT TEST III			Date/Session	11.05.2024	Marks	50	
Course code	CS3492	Course Title	Database Management System				
Regulation	2021	Duration	1:30 Hours	Academic	Year 202.	3-2024	
Year	II	Semester	IV	Departmen	t CSE		

PART A

1. What is transaction.

A transaction in the context of databases is a sequence of operations performed as a single logical unit of work. Transactions are fundamental to maintaining data consistency and integrity within a database, especially in multi-user and concurrent environments.

2. List out ACID property?

- > Atomicity
- > Consistency
- > Isolation
- ➤ Durability

3. Why is it necessary to have control of concurrent execution of transaction? how it is possible?

Controlling the concurrent execution of transactions is necessary to maintain data integrity, consistency, and isolation in a database, especially when multiple users or applications access and modify data simultaneously.

4. What are the different modes of lock?

- ➤ Shared Lock (S-Lock)
- Exclusive Lock (X-Lock)

5.Define deadlock.

Deadlock is a situation in a database where two or more transactions are waiting for each other's locked resources, causing each transaction to be indefinitely blocked and unable to proceed.

6. List the four conditions for deadlock.

The four conditions for deadlock are mutual exclusion, hold and wait, no preemption, and circular wait. All must be present for a deadlock to occur, where each process holds resources needed by the others.

7. What is meant by concurrency control?

Concurrency control is the management of concurrent transaction execution to ensure data integrity and consistency. It uses techniques like locking and timestamping to handle multiple transactions simultaneously.

8. Define two-phase locking protocol.

The two-phase locking protocol is a concurrency control method with two phases: the growing phase (acquiring all locks) and the shrinking phase (releasing locks). It ensures conflict-serializability and prevents data conflicts.

9. What is meant by scrialization? How is it tested?

Serialization is the process of ensuring that transactions execute in a sequence that produces the same result as some sequential order. Testing involves conflict-serializability and view-serializability to validate correctness

10.Differentiate strict two-phase locking and rigorous two-phase locking protocol.

In strict two-phase locking, locks are released only after a transaction completes. In rigorous two-phase locking, all locks are held until the transaction commits or aborts, providing stricter control and preventing cascading rollbacks.

PART B

11a. Discuss the violation caused by each of the following: lost update problem, dirty read, non-repeatable read, and phantoms with suitable examples.

Lost Update Problem:

- o Occurs when two transactions read and update the same data, but one update is overwritten by the
- o Example: Two users modify the same bank balance, and one update is lost due to concurrent

Dirty Read:

Happens when a transaction reads uncommitted data from another transaction.

Example: Transaction A reads data updated by Transaction B, but B later rolls back, leading to

Non-Repeatable Read:

o Occurs when a transaction reads the same data twice and gets different values due to another

Example: Transaction A reads a value, then Transaction B modifies it, causing A's second read to

Phantoms:

o Arises when a transaction re-executes a query and finds additional rows due to another

Example: Transaction A retrieves a set of rows based on a condition, while Transaction B inserts

11b. Explain Deadlock in detail with an example.

Deadlock Definition:

A deadlock is a situation where two or more transactions are waiting indefinitely for resources

Conditions for Deadlock:

- o Mutual Exclusion: Only one transaction can hold a resource at a time.
- o Hold and Wait: Transactions hold resources while waiting for others.
- No Preemption: Resources cannot be forcibly taken from transactions.

Circular Wait: A circular chain of transactions exists, each waiting for the next.

Example:

Transaction A locks Resource X and waits for Resource Y, while Transaction B locks Resource Y and waits for Resource X, causing bo h to be blocked.

12a.i) Write short notes on transaction concepts and ACID properties.

Transaction Concepts:

- Transactions are a unit of work in a database, treated as a single, indivisible process.
- Ensures data consistency and integrity through sequential and complete processing.
- o Supports operations like Commit (saving changes) and Rollback (reverting changes).

ACID Properties:

- Atomicity: Ensures all parts of a transaction are completed, or none are.
- o Consistency: Maintains data integrity before and after a transaction.
- o Isolation: Ensures transactions execute independently without interference.
- Durability: Guarantees that committed changes are permanent, even in case of failures.

12a.ii) Explain conflict serialization and view serialization.

Conflict Serialization:

- o Refers to a schedule where transactions execute in a conflict-free order, resulting in a serial order.
- o Ensures that conflicting operations (e.g., read/write on the same data) are scheduled to prevent data inconsistencies.
- Conflict-serializable schedules are tested using techniques like precedence graphs.

View Serialization:

- Ensures that a schedule has the same final output as some serial execution of transactions.
- View-equivalent schedules may have different orders but lead to the same database state.
- o Ensures data consistency without requiring transactions to execute in strict order.

12b) Define functional dependencies. How are primary keys related to FD's?

In relational database design, a functional dependency (FD) is a relationship between two attributes or sets of attributes within a relation (table). A functional dependency describes how one attribute's value is determined by another attribute or combination of attributes.

Notation: X→Y, meaning "X determines Y" or "Y is functionally dependent on X."

Types of Functional Dependencies

- Trivial Functional Dependency
- Non-Trivial Functional Dependency
- Full Functional Dependency

> Partial Functional Dependency

Primary Keys and Functional Dependencies

A **primary key** in a relation is an attribute (or a set of attributes) that uniquely identifies each row in the table. Primary keys are closely related to functional dependencies:

- Primary Key as Determinant: In a table, the primary key should determine all other attributes, meaning the primary key has a full functional dependency on every other attribute in the table.
- Minimal Superkey: A primary key is the minimal superkey. A superkey is any set of attributes that can uniquely identify a row. The primary key is the smallest subset of a superkey that maintains this unique identification.
- No Partial Dependencies: In a well-normalized table (at least in 2NF), there should be no partial dependencies on the primary key. This means that no attribute should be functionally dependent on only part of a composite primary key.

Register Number:



UNIT TEST IV			Date/Session	04.07.2024	Marks	50			
Course co	de	CS3492	Course Title	Database Management System					
Regulation		2021	Duration	1:30 Hours	Academic Y	ear 20	2023-2024		
Year		II	Semester	IV	Department		CSE		
COURSE	OUTCO	OMES	the state of the s		The state of the s		<u> </u>		
CO1:	Con	struct SQL Qu	ieries using relational a	lgebra					
CO2:	Desi	gn database us	sing ER model and norr	nalize the database					
CO3:	Cons	truct queries t	o handle transaction pro	cessing and mainta	in consistency	of the detail			
CO4:	Construct queries to handle transaction processing and maintain consistency of the database Compare and contrast various indexing strategies and apply the knowledge to tune the performance of the database								
CO5:	Appraise how advanced databases differ from Relational Databases and find a suitable database for the given requirement.								

Q.No.	Question		
	PART A	CO	ВТ
1	(Answer all the Questions $10 \times 2 = 20 \text{ Marks}$)		
•	What is B-Tree?	CO4	R
2	Define rotational latency time?	C04	K
3		CO4	R
	What is an index?	CO4	-
_4	What are the types of storage devices?	C04	U
5		CO4	U
	What is hashing file organization.	COA	 _ _
6	Differentiate static and dynamic hashing	CO4	R
7		CO4	A
	What is called query processing?		
8	Explain "Query optimization"?	CO4	R
9		CO4	A
	Define Primary and Secondary Indices?		A
10	What are the Types of Ordered Indices.	CO4	R
Service .	yr or or ordered midices.	CO4	R
	PART B	004	K
1a	Define RAID and Briefly Explain RAID		一種
	and Briefly Explain RAID	1 00 1	=
1b	Discuss two 1	CO4	E
20	Discuss (WO phase locking marks 1		_
2a	Explain about Query optimization with neat diagram	CO4	E
2b	Give detail Evel OR	CO4	$\frac{\mathbf{E}}{\mathbf{C}}$
20	OR Give detail Explanation about Hashing & Types of Hashing.	204	
	Types of Hashing.	CO4	E
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Course code CS3492 Course Title			Database Management System				
Regulation	2021	Duration	1:30 Hours	Academic Year	2023-	.2024	
Vear	ear II Semester		IV	Department	CSE	2024	

PART A

1. What is B-Tree?

A B-Tree (Balanced Tree) is a self-balancing tree data structure that maintains sorted data and allows for efficient insertion, deletion, and search operations. B-Trees are widely used in databases and file systems because they can handle large volumes of data and are optimized for systems that read and write large blocks of data (e.g., disks or other secondary storage).

2. Define rotational latency time?

Rotational Latency Time (also known as rotational delay) is the time taken for the desired sector of a spinning hard disk or other rotating storage media to rotate under the read/write head so that data can be accessed. This delay is one of the key components that impact the overall time to access data on a traditional hard drive.

3. What is an index?

An index in the context of databases is a data structure that improves the speed of data retrieval operations on a database table at the cost of additional space and maintenance overhead. Indexes are used to quickly locate rows in a table based on the values of one or more columns, improving the performance of SELECT queries.

4. What are the types of storage devices?

Storage devices include primary storage (RAM), secondary storage (HDDs, SSDs), and tertiary storage (tape drives). They vary in speed, cost, and durability, serving different data storage needs.

5. What is hashing file organization?

Hashing file organization uses hash functions to assign data to specific locations in storage, enabling fast retrieval. It's efficient for direct access but may require collision handling mechanisms.

6.Differentiate static and dynamic hashing.

Static hashing uses a fixed number of storage buckets, which may cause overflow as data grows. Dynamic hashing adjusts the number of buckets based on data volume, allowing for flexible storage expansion.

7. What is called query processing?

Query processing is the sequence of steps the database system uses to interpret and execute SQL queries. It includes parsing, optimization, and execution to retrieve results efficiently.

8.Explain "Query optimization."

Query optimization is the process of selecting the most efficient execution plan for a query, reducing response time and resource usage. Optimizers consider factors like indexes, join methods, and execution costs.

9. Define Primary and Secondary Indices.

A primary index is based on a table's primary key, ensuring unique and ordered entries. Secondary A primary index is based on a table s primary key, ensuring indices are additional indexes on non-key columns, used to improve query performance on frequently searcher

10. What are the Types of Ordered Indices?

Ordered indices include dense and sparse indices. Dense indices have entries for every record, while sparse indices only index certain records, reducing storage space at the cost of slightly slower lookups

PART B

11a. Define RAID and Briefly Explain RAID

Definition of RAID:

RAID (Redundant Array of Independent Disks) is a data storage technology that combines multiple physical drives into one unit to improve performance and reliability.

Types of RAID:

- RAID 0: Data is striped across multiple disks, increasing speed but without redundancy.
- o RAID 1: Mirrors data across disks, providing redundancy but no speed improvement.
- o RAID 5: Stripes data with parity, balancing speed and redundancy; requires at least three disks.
- o RAID 6: Similar to RAID 5 but with extra parity for increased fault tolerance; requires at least

11b. Discuss two phase locking protocol and strict two phase locking protocols?

The Two-Phase Locking Protocol (2PL) is a concurrency control protocol used in database systems to ensure serializability of transactions (i.e., the execution of transactions is equivalent to some serial execution of those transactions). The protocol ensures that the database system remains consistent and avoids issues like lost updates, temporary inconsistency, or deadlocks during concurrent transaction processing.

The Two-Phase Locking protocol operates in two distinct phases:

1. Growing Phase:

- o A transaction can acquire locks but cannot release any locks.
- The transaction can acquire any number of locks, whether shared or exclusive, as long as it does

2. Shrinking Phase:

- Once a transaction releases its first lock, it enters the shrinking phase.
- During this phase, the transaction can only release locks and cannot acquire any more locks.

'roperties of Two-Phase Locking:

- Serializable: The protocol guarantees serializability, ensuring that the execution of transactions is equivalent to some serial execution, which preserves the database's consistency.
- Deadlocks: Two-phase locking does not inherently prevent deadlocks. Deadlock detection and resolution
- Concurrency: 2PL ensures correct transaction execution at the cost of some concurrency because transactions cannot hold locks indefinitely (they must eventually release them).

2a. Explain about Query Optimization with Neat Diagram

Definition:

 Query optimization is the process of choosing the most efficient way to execute a SQL query to reduce response time and resource usage.

Steps in Query Optimization:

Parsing: Analyzes the syntax and checks for errors in the query.

o Plan Generation: Creates various execution plans to retrieve the data.

Cost Estimation: Evaluates the cost of each execution plan based on factors like I/O and CPU usage.

Plan Selection: Chooses the lowest-cost execution plan for query execution.

12b. Give Detailed Explanation about Hashing & Types of Hashing

Hashing:

 A file organization technique that provides fast data retrieval by converting keys into hash codes, which map to data storage locations.

Types of Hashing:

 Static Hashing: Uses a fixed number of buckets; collisions are handled using overflow buckets, but performance may degrade over time.

 Dynamic Hashing: Adjusts the number of buckets dynamically based on data size; helps in handling growing data sets efficiently.

 Extendible Hashing: Uses a directory that grows with data, providing a flexible structure to minimize overflow.

 Linear Hashing: Gradually adds buckets as data grows, balancing between storage utilization and access speed.

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MODEL EXAM I

Date/Session

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Course co	ode CS3492	Course Title	Database Manage			
Regulatio	on 2021	Duration	3:00 Hours	Academic Year	2023-2024	
Year	II	Semester	IV	Department	CSE	
COURSI	EOUTCOMES			•		
CO1:	Construct SOL Queri	ies using relational alg	gebra			
CO2:	Design database usin	g ER model and norn	nalize the database		- 1-1-1	
CO3:	Construct queries to	handle transaction pro	cessing and maintain	in consistency of the	e database	
CO4:	Compare and contra	st various indexing st	rategies and apply th	ie knowledge to tur	ie the periorii	
CO5:	Appraise how advance	ced databases differ fi	rom Relational Data	bases and find a sui	table database	e for
	the given requiremen					
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* 2. 11			PART A	E-1844 4 F		
		(Answer all the Q	uestions $10 \times 2 = 20 \text{ Ma}$	irks)	V71	
1	Define data abstraction.	20			CO1	R
2	What is full outer join.				COI	R
3	Short note on second no	CO2	A			
4	List the feature of EER	CO2	A			
5	Define deadlock.	CO3	R			
6	List the four condition	for dead lock.			CO3	U
7	What are the states of t	ransaction?			CO6	R
8	Define two phase locki	ng protocol.			CO6	A
9	When is a transaction r				CO4	R
10	Define upgrade and do	wngrade.			CO5	R
No. ar ar	2 3 1 2	190 - 20 20	PART B	ATT - Tel 3 - (2.4 g b.)	Market S	11.00
		(Answer all the Q	uestions $5 \times 13 = 65 \text{ Ma}$	arks)		
110	Explain DBMS system			The same State of the same of	COI	E
11a	Explain DBIVIS System	arcintecture				
			OR	N)	2 1	
11b	Explain about relational	model constraints			CO1	E
110	S.p.a accur			r ē	1	
12a	Draw E-R diagram for l	nospital managment.			CO2	C
			OR			
12b	Explain briefly about n	ormalization		9	CO2	E
13a	Explain in detail abou	t ACID property.		1 2	CO3	E
-			OR			

135	Explain log based recovery in detail	CO3	C
	i) Immediate database modification		
il S	ii) Deferred modification		
14n	Explain all types of data models.	CO4	A
	OR		
146	Explain various DML command with neat syntax.	CO4	E
15a	What is concurrency control? How is it implementation in DBMS?	CO5	A
7.5%	OR		
15b	Explain E-R model concept and extended E-R model	CO5	C
	PART C		
	(Answer all the Questions 1 x 15 = 15 Marks)		
l fia	Consider the employee database, where the primary keys are underlined Employee (empname, street, city)	CO6	C
	Works(empname, companyname, salary) Company(companyname, city)		
	Manages(emphame, management)		
	Give an expression in the relational algebra for each request		
1	1. Find the names of all employees who work for first hank company?		
-	2. The the halles, street address and cities of residence of all		
-	and earn more than 2000000 more annual	a y	
1	3. Find the names of all employees in this database who lives in the same city as the company for which they work.		
1	4. Find the name of all employees who earn more then every employees of smallbank		
	corporation corporation		
	OR		
bI	Describe briefly about embedded SQL		

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	MODEL EX	KAM I	Date/Session	01.06.2014		
Course code	CS3492	Course Title		01.06.2024	Marks	100
Regulation	2021	Duration	Database Management System			
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		Semester	IV	Department		CSE

PART A

1.Define Data Abstraction.

Data abstraction in DBMS hides lower-level details to simplify data access for users. It provides three levels: physical (storage details), logical (structure and relationships), and view (user-specific perspectives). This helps in efficient data management.

2. What is Full Outer Join?

A full outer join returns all rows from both tables, with NULLs in columns where there is no match. It ensures that unmatched rows from both sides are included in the result.

3. Short note on Second Normal Form.

A table is in Second Normal Form (2NF) if it is in First Normal Form (1NF) and all non-key attributes depend fully on the primary key. This eliminates partial dependency and ensures better data organization.

4. List the features of the EER Model.

The Extended Entity-Relationship (EER) model enhances ER models by adding subclasses, superclasses, specialization, and generalization, helping to create more complex and detailed database designs.

5.Define Deadlock.

A deadlock is a situation where two or more transactions are waiting indefinitely for each other to release resources. This results in a standstill, where no transaction can proceed. List the four conditions for deadlock. Deadlock occurs under four conditions: mutual exclusion (one resource per transaction), hold and wait (holding resources while waiting for others), no preemption, and circular wait.

6. List the four conditions for deadlock.

Deadlock occurs under four conditions: mutual exclusion (one resource per transaction), hold and wait (holding resources while waiting for others), no preemption, and circular wait.

7. What are the states of a transaction?

A transaction has multiple states: active (executing), partially committed, committed (successfully completed), failed, and aborted (rolled back due to failure).

8.Define Two-Phase Locking Protocol.

The two-phase locking protocol ensures transaction serializability by dividing the process into two phases: a growing phase (acquires locks) and a shrinking phase (releases locks). This helps prevent conflicts.

9. When is a transaction rolled back?

A transaction is rolled back when it cannot complete successfully, such as during errors or deadlocks, Rolling back restores the database to a consistent state.

10.Define Upgrade and Downgrade.
In DBMS locking, upgrading changes a read lock to a write lock for higher access, while downgrading reduces a write lock to a read lock, adjusting access levels based on transaction needs.

PART B

11a. Explain DBMS System Architecture

- DBMS architecture includes three core levels: internal, conceptual, and external, which work together for
 effective data management.
- Internal level: Manages physical storage details, including data compression, indexing, and data structure. It deals with efficient storage and retrieval mechanisms.
- Conceptual level: Provides an abstract structure for the entire database, independent of storage and focused on logical organization.
- External level: Presents specific views to users, showing only the relevant information, enhancing user experience and data security.
- This layered structure supports data independence, efficient access, and data abstraction in complex systems.

11b. Explain Relational Model Constraints

- Relational model constraints maintain database integrity and accuracy, enforcing rules on data.
- Key Constraints: Primary and foreign keys ensure unique identification of records and maintain relationships between tables.
- **Domain Constraints**: Restrict data types for each attribute to ensure correct data entry, like integer-only for age fields.
- Referential Integrity: Maintains consistent relationships by ensuring foreign key values exist in referenced primary keys.
- Entity Integrity: Requires primary keys to be unique and not null, preventing duplicates and ensuring every record is identifiable.

12a. Draw E-R Diagram for Hospital Management

- The E-R diagram for a hospital system includes entities like Patient, Doctor, Nurse, Appointment, Treatment, and Billing.
- Attributes: Each entity has unique attributes, such as Patient (ID, name, address), Doctor (ID, specialization), etc., to capture detailed information.
- Relationships: Define associations, such as Patient and Doctor through Appointment, or Doctor with Treatment, to reflect interactions within the hospital.

• Cardinality: Specifies one-to-many or many-to-many relationships, like a Patient having multiple Appointments, helping manage hospital operations effectively.

12b. Explain Normalization Briefly

- Normalization is the process of organizing data to reduce redundancy and enhance integrity by structuring tables and dependencies.
- 1NF (First Normal Form): Ensures atomic values in each field and removes repeating groups or duplicate columns.
- 2NF (Second Normal Form): Builds on 1NF by eliminating partial dependencies, ensuring non-key attributes depend fully on the primary key.
- 3NF (Third Normal Form): Removes transitive dependencies, meaning non-key attributes are independent of each other, leading to better data integrity.
- BCNF (Boyce-Codd Normal Form): A higher level where every determinant must be a candidate key, further refining the database structure.

13a. Explain ACID Properties in Detail

- ACID properties are critical for reliable transaction processing in databases.
- Atomicity: Ensures a transaction is completed fully or not at all, with rollbacks on failure to maintain data consistency.
- Consistency: Guarantees that database rules are enforced at all stages, so the database transitions from one valid state to another.
- Isolation: Prevents transactions from affecting each other's execution, which is essential for accurate
 concurrent transactions.
- Durability: Ensures that once a transaction is committed, it is permanently recorded, even if the system
 crashes.

13b. Explain Log-Based Recovery in Detail

(i) Immediate database modification

Immediate Database Modification

Immediate Database Modification is a recovery technique where changes made by a transaction to the database are written immediately to the database before the transaction commits. These changes are then recorded in the log, and the log entries are used to determine the recovery actions in the event of a failure.

Key Concepts in Immediate Database Modification:

Write-Ahead Logging (WAL): In this approach, the log is written to disk before the actual data is
written to the database. This ensures that if a crash occurs, the system can use the log to recover and roll

back any uncommitted changes. Specifically, any changes to the database must be preceded by a log entry.

2. Transaction Commit: When a transaction commits, the database has already been updated (because changes are written immediately). The system writes a commit record to the log to indicate that the transaction has completed successfully. This record tells the recovery system that the transaction should be considered permanent.

3. Crash Recovery: After a system crash, the recovery process will use the log to determine which transactions were committed, which were rolled back, and which ones were in-progress at the time of the crash. It uses the following steps:

- o Redo: Reapplies the changes from committed transactions to ensure that all committed updates are reflected in the database.
- o Undo: Reverts the changes made by transactions that did not commit before the crash.

ii) Deferred modification

Deferred Database Modification is a log-based recovery technique where changes made by a transaction are not written to the database until the transaction has committed. Instead, the transaction's operations are logged in the transaction log, but the actual changes are deferred (i.e., not applied to the database) until the transaction is committed.

This method relies on write-ahead logging and ensures that only committed transactions can modify the database. If a system crash occurs before a transaction commits, the changes made by that transaction are discarded because they have not yet been written to the database.

14a. Explain All Types of Data Models

- Various data models structure databases according to different needs and relationships:
 - o Hierarchical Model: Organizes data in a tree format with parent-child relationships, suitable for hierarchical data.
 - o Network Model: Extends hierarchical relationships to many-to-many, forming a graph structure for complex associations.
 - o Relational Model: Uses tables (relations) to represent data, focusing on relationships between tables through keys.
 - o Object-Oriented Model: Stores data as objects, like in programming, allowing for complex data types and relationships.
 - o ER Model: Visualizes data with entities and relationships, often used as a design framework before building the database.

14b. Explain Various DML Commands with Syntax

- Data Manipulation Language (DML) commands allow manipulation of data within tables.
 - SELECT: Retrieves specific data from tables, with options for filtering and joining (SELECF * FROM table_name WHERE condition).

- INSERT: Adds new records to a table with specified values (INSERT INTO table_name VALUES(...)).
- UPDATE: Modifies existing records, typically with a condition to target specific rows (UPDATE table_name SET column=value WHERE condition).
- DELETE: Removes records from a table based on specified conditions (DELET) FROM table_name
 WHERE condition).
- MERGE: Combines INSERT and UPDATE operations based on matching conditions, providing a flexible data modification option.

15a. What is Concurrency Control? How is it Implemented in DBMS?

- Concurrency control manages the simultaneous execution of multiple transactions, preserving database consistency.
- Locking Mechanisms: Use shared and exclusive locks to prevent conflicting access to the same data by multiple transactions.
- Timestamp Ordering: Assigns timestamps to transactions, enforcing an order to avoid conflicts in execution.
- Two-Phase Locking Protocol: Divides transactions into growing (acquiring locks) and shrinking (releasing locks) phases to avoid deadlocks.
- Optimistic Concurrency Control: Assumes minimal conflict, checking for issues only at commit time to improve transaction speed.

15b. Explain E-R Model Concept and Extended E-R Model

- The E-R (Entity-Relationship) Model represents data with entities, attributes, and relationships, providing a framework for data organization.
 - Entities: Represent real-world objects, like Student, Course, or Instructor, with unique attributes such as name, ID, or course title.
 - Attributes: Define properties of entities, such as name, age, or address, to capture detailed information about each entity.
 - Relationships: Represent associations, like a Student "enrolls in" a Course, specifying how entities interact with each other.
- Extended E-R (EER) Model: Adds complexity with subclasses, superclasses, specialization, and generalization, allowing for inheritance and hierarchy in data relationships.

16a. Consider the employee database and give relational algebra expressions for each request

- Employee Database Structure:
 - 1. Tables include:
 - Employee: (empname, street, city)
 - Works: (empname, companyname, salary)
 - Company: (companyname, city)
 - Manages: (empname, management)
- Relational Algebra Expressions:
 - 1. Find names of all employees working for First Bank Corporation:
 - Use selection and projection to identify employees associated with "First Bank Corporation."
 - Expression: π_empname(σ_companyname='First Bank Corporation'(Works))
 - 2. Find names, street address, and cities of employees at First Bank Corporation with salary > 200,000:
 - Use selection and join on Works and Employee to filter for company and salary conditions.
 - 3. Expression: π _emphase, street, city(σ _companyname='First Bank Corporation' \wedge salary > 200000
 - 4. Find names of employees living in the same city as their company:
 - Use selection and join on Employee and Company to match city fields.
 - Expression: π_empname(σ_Employee.city = Company.city
 - 5. Find names of employees earning more than every employee at SmallBank Corporation:
 - Use division or nested query approach to compare salaries across companies.
 - Expression: π_empname(Works) π_empname(σ_salary < ALL(π_salary(σ_companyname = 'SmallBank Corporation' (Works))))

16b. Describe Embedded SQL Briefly

- Embedded SQL:
 - Definition: Embedded SQL is SQL code embedded within a programming language like C, Java, or Python. It allows applications to interact with the database directly within the code, simplifying database operations.
- Components of Embedded SQL:
 - o **SQL** Statements: Standard SQL statements, like SELECT, INSERT, UPDATE, and DELETE, are used to query or modify data directly within application code.
 - o Host Variables: Variables in the host programming language can be used to pass values into SQL queries, making them dynamic and adaptable.
 - SQL Control Statements: Use constructs like cursors and error-handling statements to manage database interactions efficiently.
- Process:
 - o **Pre-compilation**: The embedded SQL code is pre-compiled to ensure compatibility with the database, creating efficient access paths.
 - Execution: When the application runs, SQL statements are executed as part of the program, allowing real-time data processing within the application.
- Advantages:

 Efficiency: Reduces the overhead of separate database calls by embedding SQL, improving performance.

Data Integrity: Enforces data checks directly within the application, helping to maintain data

consistency.

Seamless Integration: Allows developers to write SQL queries in the application's programming language, bridging database and application code efficiently.



Answer Book

Name	V. Hari	วขึ้นล		Year/ Semester/Secti	on [] IV
Register Number	732422104	/	27.3.24	Department	CSE
Course code	CS3492	Course Title	Databas	e Managemer	nt System
Internal Assess	sment Test	IAT 1	IAT 2		Model
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b. Data Definition Language: 200 a expecial substanguage who a expecial substanguage who a expecial jacustates to access the datas. Su that is a language which is provided creating and modiffying the table, view, underson on was account of the stable, view, Common command for Dollare * CREATE O JIN SA * ALTER 10 Data Base System JOAR CHAROLD * Data reductioned * Data reductioned 7. NULL Valuein So *. In Database NULL value provides special value to pata base. They are en two cases (P) when fletal name of some tuples are unknown (ine) city, name is not given (B) inapplicable lies middle name & not present wandards distruct * There B. about * more Ps 20 8. key: of the discussion of keys are * Super Koy * Fornegn Key * Primary Koy * Candidate Key.

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PART-B. TO feter the students have more 18. Relational Algebra: 81 200 MONE Relational talgebra us a query language used to specify the relation to read the table of different ways. ARRE Relational operation 988 Projection no Cartesian selection product olded (Relational of monutar ration operation Set Join) operation operation Natural Pour Conditional) J (Face) 10 1 Cunton operation John unitorsection Student Selection: INA JMANS 013 selection operation of used to getch tuples con a relation (table) Syntax: Opredicate (relation) Example: Student (relation) age SID sName day the column 1800 HARA 001 BBBB 23hld6 1000 003 16 CCCC

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STAHS. To fotch the students have more Query: than age 18. assispité desirables. Visus Jages 18 Cstudient) Lonaito 1951 and butput sname 1 9113 PAAA BBBB matherene levels 199 DODI Phofe elitor -Selection Propertion: Projection & used to project some columns en la Rolation (table) Thus to alloplay particular column in a Rolation. Dpody Syntax: But TICI, Co. .. Welation distance Example: student SID SNAME Selection 3NA book 101 AAA selection.L colotion (table BBB 03 CCC OLL Example: milled facery: ofa same sig wei wantito desplay the columns et student name and age only. TI (ISHAMIE) AUE) (Student)

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output: John over Linn SNAME AUE ARROYED to comition 1800 Bull Coldato Babos 835 a 19 gwd use chave in elatins. 19.8 turco ont DDD 88 ataka Hodaga CartesPan Product: 30% Londitioned (2) Cartespan product by wed to combine two differentia vielations brable du Some widdling syntax. Syntax: 9,MA Example: student Roewenex SID TUBM DAY CNAME ALLE OID. 20 0 00 m 100 5 2 8 18 19 POOL SARAIDIS ०१८१८ ० ८०० व नाश्रीव 0 19 TOBBB 0102 03 3 5006 009/8/19 080 CCCC 103 osletos + 5001 40000 181 OH 18 405 40 FOO 140 વલા ના દા Query: want to combine who save renerve Poblaz 005 Julion and stradent, and reuse nelation under (Studential & ROOM & SIDD A) (RODENE SISTINGOUS) (student x Roserve) Student 1. Igerondent and post of Inghis PAE ISBM PAY dame SID 005 8/8/19 18 AAA 01

-1/8/19

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Soin operation Burbed to combine two or more relation (tables) then the combined relation we chave to yetch data. 55

(P) conditional jon : toubor! distration solding of the operation & wood to comber two relations (tables) under w some conditions

Syntax.

A Mcl

Example

Sypotax

Example

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14	STOOD .	W.		

(ii) Equi join:

* It is a kind of join equality condition between relation (table) we equalate the relation is a oque foir.

unsbord

(iii) Natural Joba:

*. There are common columns in a relation when we want to combine the common columnicity using natural Join of the matural foundition in or imple (M). WHUS

managa 11. DataBase System Architecture:

* Data Base system Based on relational idata model

*. It shows the application interpre upos narve uper, application program created by application programmer, Query dook are used by sophisticate user and administration took were used by patabose ad mb toution

(support Bullitz motors special)

odabane SophPlate Appliation administration Narve user user WIR une 12/willen admenstate Query Application KOL Application toob. took program Phterfaces compiler & DMLC Pritorpreter Jugaker complied DMI complies VI program and dosigen. cobject or 9191 4 code Query evaluation Processor Authorization of transation Buffer 4 Portegrity manage al Aldrorage derice 010 Docts Louandles willingtoxage Duta Prolices Coata Base System Archi tecture

DATE

*. The lower level architecture B storage albk. * storage device us wentprevet conjust There two main component of data base system soft bus south * Query processor donce at yotorages device. squalbit Query processor: It include the * The iquery processor unteractive one used to specify of the a specification of the specification to access the dottain off it * It uncludes a DOL contempreter, DMI compler and Query ovalution * Date andra ungeno. * Date (Detucnony: DDI autorbiotos; *- It & a translator in which it transmits query en data dictionary DMI compiler: *It Ba translator which transmit the data anto overy englis plan. Overy evalution ergine.

* It executes the dow doubt unstructions your the DY2 comples

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Answer Book

Name	P.Res	athi		Year/ Semester/Section	n I iû
Register Number	7324 22104	Date/Session	17.4.24	Department	CSE
Course code	CS3492	Course Title	Databa	se Management	System
Internal Assessment Test IAT 1				IAT 3 Mode	V
Name and Sign	nature of the Invigi	lator with date			

Instruct	Instruction to the Student: Put tick mark to the question attended in the column against question.									
1	Part	A	Part B/ Part C							
Q. No.	~	Marks	Q. NO.	V	a	V	b	Total Marks		
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Course Outcomes	1	2	3	4	5	6	Total
Marks allotted	50	50	_	^	-		50
Marks Obtained	44	16	_	_		2	46
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Unit Lest - IP. Dame: P Revorthi 900 10 8 732422104086 Dept: BE-CSE/II year.
Subject: Data Base Management Syste
Subject Code: CS:3492. Data: 17/04/2024. bell-A. A. Hard March of 100 Explain Entity Relationship Model: the Entity Pelationship b donoted who by BRA Eddins of Date to represent the Entroprises of Schema up the Date of Structions of Logical and database the ER Model is Represent the Schema of Entity Relationship. 2) Give the limitation of B-R Model? How do out you Over Come +6899 and to (1) Loss of Information Content. the Loss of the entity from the ERMod (2) dess time of representation of Ex Model. (3). No representation of Maniputating Tator - they are Manipulated from the ER Model.

Jugh Level denge - they Can be high devel of Data wasterign unto 3) What is an entity: -. the Data Base us with distinguisted to Some officer Objects there is Called Entity. For Example: They Student name of "poonain" the entity of the open to Student Mame can be 4) Défine functional Dependancy: * the two column of Sot us pard a stru dable of functional popendo unitur of the p-> a then they have two row of the Dependent of entitle are equal on the people on the fools. TI.P=T2.P and TID=T2.D they agree refunctional Dependant with The of perojet is (T) and autibute
of the (P) we therefunctional Dependary

5) touval fundfoual Dependency:got of A 18 the value of Dependancy
B 18 the MVD of the functional
Dependancy 18 the A1B 18 the
burral of B. then the B 18 the equalent values. For example: - the table two Column Mames and A and B the inlight there is a faulfal order of Dependancy In the raime, cause i'd is Calculud b) Define Mormalization:-. * The normalization is reorganisate of the Data to need from the Data * Huy are Conditions of Data Pr normalization. * No dependancy of the table * Data redundancy of the table * It is imposit of Database takeup Renformense. Performense.

7) Kultfralued Depandancy? It the table is should have Multivalued Dopendancy it Sally the following Condition, town * It should set of colour is A>B B the A/B the fenctionally Dependency, B & the Multivalued Dependancy there table 12 valued dependancy. It should table has at -Lea cloudons us the rultiva they have Relation of RCABC) A and B, then Band C is the Inderpendent of the table is the MultPralued Dependancy. BCNOF and the relation which we BCNE? BONF IN Deproted by the Boycecodd Normal For. of the relation with the Certain is the BCN F. non relation of Dependany

Offilm

It can be third pormal form of BCNF should satisfy the form. Indomedia del de mios il 18 BONE anahal is not is 3NR At there is called BCNF. Cource Reva chennai. Dharstn. BB necessary to Decompose a relation: they are norsary to the Decomposi a Julation is one Strigle Lable Value break down by the Multi value us Decompose. 1- to Wantowillograpist. * Here is Decompose of the relation us olumnated Porto the relation Two Destrable properties of Delonipositions. (1) Loss-Loss or now Less 2. It is have the Decompose J- Atroved tothe original Duta Storage Date Base

A Improvement.

2) Depondant revention- Et Bhould Information can handlet Property of the chelt and Paut-B.

11) Explain un Detail about Mormal For Mormal Form: * the normalization is the process of vuorganization of Data to need for the Data Base / B of planting two Paric Components . It has no Dependancy of the table A It shout have the Lable in Mormal form is Proportant of Data Bancostake up Less of the time * It also need to help for the Improved performance. the Data Base.

Ringle authbute 1 columns.

Data Base.

has unique noang

* the Leple of Drose not a matter.

several bas tasking

Example:-

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19) Multivalued Dopendanus and Fourth

rullivalued Dopendency:

Dependency is the Golfsty the following

the A B B the A B the single value Dependancy and B B the Lable is Multivalued Dependancy they are the table is Multivalued Dependency.

Columns as the table of Multivalued Dependency.

forms in Aand B then Band C is Independance of the for is MVD.

they are Satisfy the three Condition is

A B the Multivalue Dependency
B B the Multivalue Dependency of
the Huy those Columns are Independent

& It should have Multivalued Dependay
as denoted by

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* they name A'curble B is the mot have MVD is the table wishte fourth normal form.

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the table Should have stependent of the Coleon for splits of the Dates Base Normaly forms

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Fowith Cornal Form:

the following it salisty the two Condition

Mormal form

Dependency.

Example.

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Answer Book Year/ Semester/Section 11/10 Name Register 24.5.24 Department CSE Number Database Management System Course Title Course code IAT 4 Model JAT 2 IAT 1 Internal Assessment Test Name and Signature of the Invigilator with date

Instruct	ion to	the Student	Put tick ma	rk to t	he question a	ttende	d in the colum	n against question.
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		rough	(eOI			Name and of the IQA	Signature C member

132422 toxo18
BECSE PART-A MARTIN 1 B-tree: * B-tree is a sey-balancing search tree in which data is stored in Strted order and allows searches, sequential acces; insertion and deletion in logarthmie time * It's used in database and file systems to ensure that the tree remains balanced with its height Kept low for efficient operation. a Rotational latency 19me: caused by the time It takes for the desured sector of a disk to rotate under the read! write head. An Index Badatan Structure that Amproves the speed of data retornal operation. * It proxides quick acres to nouse the a database table by

4. Types of storage device: storage * Belmany * decondary storage * tertiary Storage 5 Hashing gile Organization? & Hashing file organization hach function is applied to a dey deld to determine the address of the data socon Static Hashing & ged using fechagins wike yong or open addiening. The harh table Sto can grow of Shrenk dynamically based on data charles a linear trashing

7. Query oprocoving 8 translating a high - level query written in a query language anto clow- cover instruction that the database can execute. 8. Query Optimization:

* It is the process of

Selecting the most officialt

way to exocute a given

Query.

The database management

System exaluates different execution

plans and chooses the one with the dowest costato dodon is yell 9. Paimary Index:
An Tindex on a primary
Key, which ensues uniqueness and B
directly that to the physical Secondary Index: tegs, allowing efficient access to data based on fields other than the primary key the source of the primary key the source of the source o and stell of course production of the the to dily destine.

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10. Types of ordered Indices: Dense Phodex * charse Pindex 1910 * clustroied Prodex Non-clustered Index. server acounted PART-B 11(a) dek (RAID) * originally redundant away of unexponsitio asks. is a way storing. the game data in different places on mutiple hard dok to protect data in the case of a deino failure. manage as vage number of disk, providing a view of a spingle disk of chigh capacity and high speed by ching muttiple disk in paracol and high reliability by storage data redundantly U Benefits of RAID + Data Coso can be very dangerous for an organization U
* RAID technology prevents data loss due to ask gottine.

DATE /

1 1 704 4 11 (1) RAID O: * RAID Level o divides block units and workes them across a humber of walks of any ?. in Suth # AS data B placed across multiple ask. It is also colled-data strepping
There is no parity checking
of data then so at data in
one diduce get corrupted then all
the data would be lost of in season on a survey had * Ilo performance & greatly Proproved by spreading the I/o Load
acrows many channels of dulyon.

A Best performance & achieved
when data is stropad across multiple
controlled with any only one dulyou per controller. DB advantages: * It & not gault-tolevent tailure & one dade: com societ en all data en an away being last in RAID O.

RAID Lovel 1: Su lead withilly a forced titel to # Also known as disk measured thus configuration consists of at a daylicate dast two delves that deplicate the storage of data. There is no stripong on a star strip since ofther dole an be read at the same time. waite performance is the game as for orgee dek storage. * Data loss would occur only ask also fails before the System is repailed probability of combined ovent at vory amaly. /io/lostrai RAIDI (block) Co-thurs block) Was dun Horbs blodc2 block3 dock3 CV. dock4 block 4 duno dule - Educati Late. Chub . Walls In Farity

RAID Jevel 2:

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RAID 3 RAPOL

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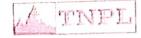
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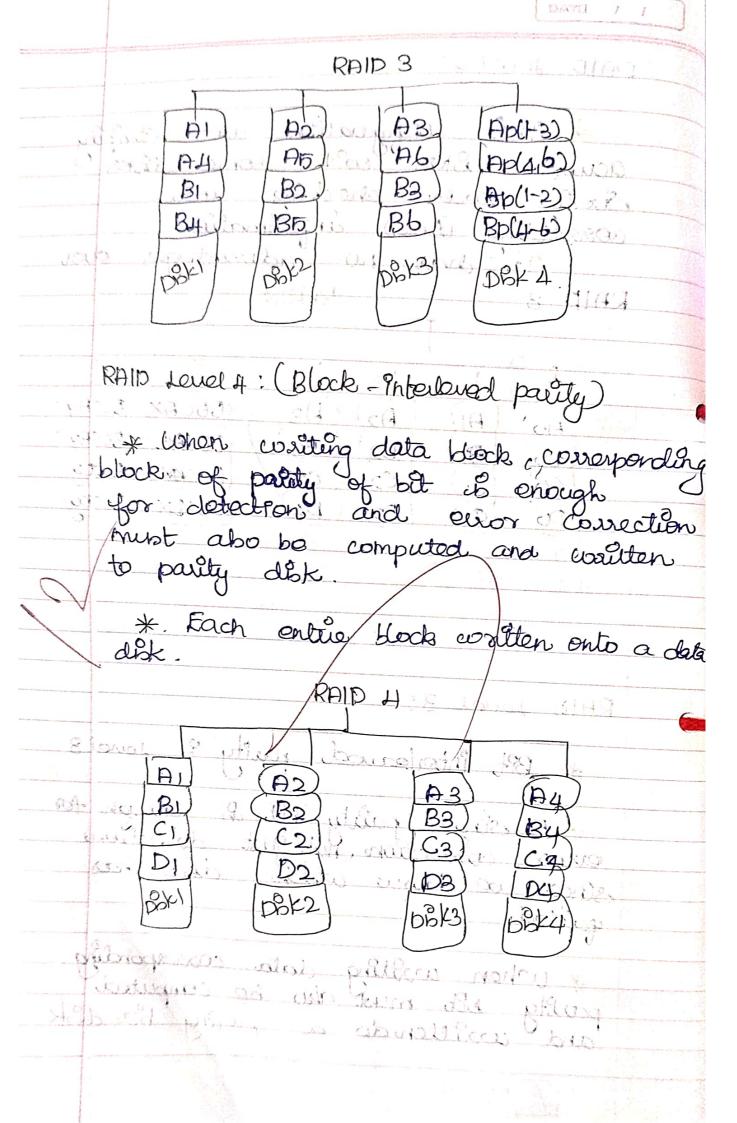
RAID Level 3: 1 11

* Bit-interleaned parity & Lovel 3

* A single parity bit is remorgh for error objection not just detection series we know which disk has golded.

* when worting data, corresponding party bit must also be computed and wortlendo a party bit disk





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(b) Hashing:

* It B a technique wed in database management aystem to officiontly retrieve data used wared on a apears areach key.

database sequentially harring uses a chain function to map a search bey to a specific docation of the memory or storage.

a constant time complexity O(1)

Types of Hasing:

+ Hashing can be broadly clarified

Static Harhing:

* In static hashing the hour function and number of buckoto are gived at the time of database cleation. This means the structure dynamically does not grow or shrink dynamically

original de par aprenenta y sur de sonos de de s

ATNEL

Advantages: * Simple to Implement

* Suitable for systems where the

as it cannot handle originate growth. Dhadwantage ".

* collision may concione as the databale gebos

collistan Rosolution Technique.

* chaining *- Open addersing

1) Jungar probbing

2) Quadratic probing

3) Doube hashing

Dyramic Harbing:

Stre of the charm table dynamically as the database grows or strucks.

7hB B consmonly used an database where the number of seconds charges frequently.

Advantages:

* Effecient database

for growing or shrinking

* Reduce overflow insues by dynamically adjunting the structures.

Methods in Dynamic Harring:

2) Linear harring

Application of Hashing?

D Indexing & Harring & widely ned for creating harr-based

caching: Hash function are used in chaning mechanisms to map data to memory docation

3) Password storage: Hashing was security by storing harried panwood

4) Data deduplication

A PERTURE



CONTENT BEYOND THE SYLLABUS

1. Advanced SQL Techniques

Window Functions: Used for complex data analysis within partitions of data.

Recursive Queries: Useful for hierarchical data, such as organization charts or folder

 Common Table Expressions (CTEs): For creating temporary result sets to use within a SELECT, INSERT, UPDATE, or DELETE statement.

 JSON and XML Data Handling: Storing and querying semi-structured data within SQL databases.

2. Indexing and Optimization Techniques

- Advanced Indexing Types: Such as partial, clustered, and composite indexes, and when to use each.
- · Query Optimization Techniques: Deep dive into the query optimizer, execution plans, and using hints to influence the optimizer.
- Database Statistics: Understanding how databases gather statistics and how it impacts query performance.
- Materialized Views: Creating precomputed views that improve performance for frequently accessed data.

3. Concurrency Control and Transaction Management

- Isolation Levels and Consistency Models: Including Serializable, Read Committed, Read Uncommitted, and Repeatable Read.
- · Optimistic vs. Pessimistic Locking: Different strategies for handling simultaneous
- Multiversion Concurrency Control (MVCC): Used in databases like PostgreSQL and Oracle to handle concurrent transactions.
- Deadlock Detection and Resolution: Methods for managing deadlocks in database transactions.

4. Distributed Databases and Sharding

- CAP Theorem: Understanding the trade-offs between consistency, availability, and partition tolerance.
- Data Partitioning: Horizontal vs. vertical partitioning, and hash-based, range-based, and composite sharding.
- Replication: Types of replication (master-slave, master-master) and their trade-offs.
- Eventual Consistency: Used in distributed systems to provide availability and partition tolerance.

5. NoSQL and NewSQL Databases

 Types of NoSQL Databases: Key-value, document-based, column-family, and graph databases, with examples like MongoDB, Cassandra, and Neo4j. • Event Sourcing and CQRS (Command Query Responsibility Segregation): Techniques

 NewSQL Databases: A newer class that provides SQL capabilities with the scalability of NoSQL, like CockroachDB and Google Spanner,

6. Data Warehousing and ETL (Extract, Transform, Load)

Star Schema and Snowflake Schema: Data modeling for analytical databases. ETL Pipelines: Tools and processes for extracting, transforming, and loading data,

including Apache NiFi, Talend, and Informatica.

 Columnar Storage: Used in databases like Amazon Redshift and Google BigQuery for analytical processing.

Data Lake vs. Data Warehouse: Understanding the differences and when to use each.

7. Database Security and Compliance

Encryption Techniques: Data-at-rest and data-in-transit encryption.

 Role-Based Access Control (RBAC) and Attribute-Based Access Control (ABAC); Advanced security models.

SQL Injection and Prevention: How to prevent injection attacks.

Compliance Standards: GDPR, HIPAA, and their implications on data storage and processing.

8. Database Backup, Recovery, and High Availability

Point-In-Time Recovery (PITR): Techniques for restoring data to a specific time.

High Availability Strategies: Failover clustering, active-active vs. active-passive configurations.

• Database Snapshots: Using snapshots for backup and quick recovery.

 Log Shipping and Replication: Techniques to ensure data is mirrored in real-time or nearreal-time.



DEPARTMENT OF COMPUTER SCIENCE AND ENGINEERING

Assignment Question Paper

	Assignment –	1	Date of Issue:	15/02/2026	Marks	10
Course code	CS3492	Course Title	Data Bare	managemer	it sy	stem
Year	l)	Semester/Section	Ñ	Date of Submission	on: 281c	3/2024

Q.No		Questions	CO
L	Explain in detail a	bout relational algebra	Col
	Mith example	0	

Name and Signature of the Faculty Incharge

HoD/CSE



DEPARTMENT OF COMPUTER SCIENCE AND ENGINEERING

Assignment Answer Sheet

Name of the Student: & one Pray

AU Register Number: 732422104043

Assignment - 1

Course code | CS31492 | Course Title | Data Brue | Vonagement System

Year | 1 | Semester/Section | V | Date of Submission: 28 | 03 | 2024

Q.No	Questions	CO
1	Explain in detail about Polational Algebra	Co1
2	With example	
3		

Mark Allocation

Marks Allocated	Marks obtained
6	5
2	1
2)
10	7
	6 2 2 10

Name and Signature of the Faculty Incharge

HoD/CSE

Assignment - I

Manue

: d. dri Raj

Dutsject

: Data Base Management systems

Rey NO

: 732422104043

Date

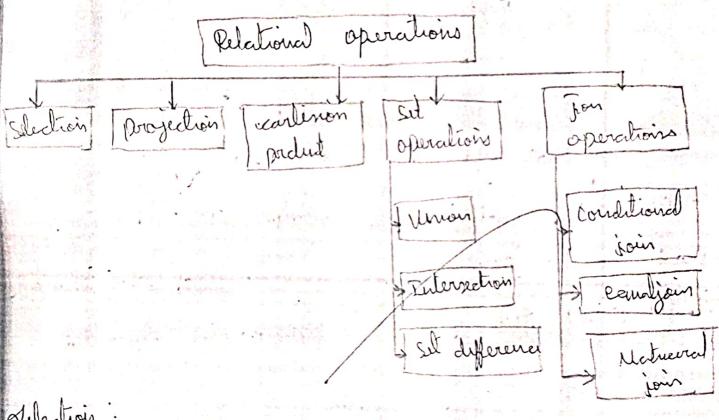
: 72.03.2024:

Ja 32 m.

Relational Algebra:

procedural serving # Relationed ralgebra is a language which is used to lacers the database tables to read data in different warrans.

the valational ralgebra # The queries present in cour denoted usuis operators.



delectrois:

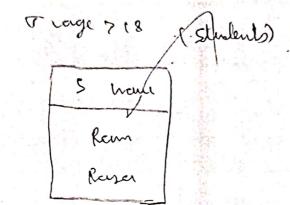
operation is used to fetch vious or This

Opredicate (table name or vulation) Syntax:

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	Rum)1	M
2	Rasa	32	M

cou:

H students with carge more than es



Pojection:

projection operation is used to project only a sentari set of realthbuts of violation. That nears if you mand to be seen only the name call of the students in the table thus you can use project aperation.

dyntax: TI C, Cz ... (Relation)

Revery: Display the name land large of all the st

and the same of the same	icicle
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Shyron	(8
seeka &	10
geeta	13

3) Cartisian product:

This is used to compile data from two different orderion (takks) into any would or fetch data from the companied valutions.

dyntar: ARB

For Example: Suppose they care two tables estimients

14.2	stulent	٥
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2	styrum	13
3	suta	16
4	Gleiter	. 18

lezeur	
TSBN	class
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007	29.03.41
007	ા.વા.બ
	005

Quay: find the nears will the students to have men ISBN = DOT to welly to Runy.

Etndent. Sid : Reserve Sid, Reserve :1500 : 001

aut put:

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5	Cherry	cereji	sid	7123 V	clay
	Clayen	13	/2	700	29.03.23
3	sute	160	5	0.2	
		Millian		700	17.09.24
	1				

det operation:

Various Set of operations ware min intern rand set difference.

union: This operation is used to fitch data from I relation (table).

Syntax: AVB

E0 .			The second second		
Example:	Student		kyst -	Book	
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1	aaq	178	000	BMS	XXX
2	bbb	18	006	05	VYX
	1	Section of the sectio			1

we want to diplay both the structure went edded all alted mort eman

Fenom (student) U output. brand Sweet Sweet DBMS aca bibbb This apprehation is used to filets/data-from both the Intersection : tables calich vio common en the both tables Syntex: ANB Exempli: student Branch Name Salary 50000 ana 届 aaa CSE ववीव 60000 bbb EJ Query ! If we went to find out the homes of the shudents who core working in a company. To vane (student) A Tome (worker) Out put: Name ppp

The cresult of sit difference operation is duples, but Set - ifference early present in an endation in the sent took in Grand . (while) waterline Symlax: A-B Excupite: - consider due vulation full him cand part low employer, if we want to out cooking for full time Han find the set -differency. Il emphane (Eul) . time. employell) -Tempram (part - trii - emplanges). Join operation: * A join operation is used to common information from two or more valation. * Join aparation is used cas M. Conditional Join: This is can aperation information from town to

AM.B = J. (AXB)

Equal join:

This is a kind of join operation in which I

ratural join:

ratural join:

ratural country common common country with the country common columns can use the control of control of control of columns.

Synday : A MB.



DEPARTMENT OF COMPUTER SCIENCE AND ENGINEERING

Assignment Answer Sheet

Name of the Student: B. Saxam

AU Register Number: 732422 104038

Assignment - 2

Date of Issue: 15/04/2024 Marks 10 Course code CS3492 Course Title DataBase Year Semester/Section

Q.No	Questions	CO
ا عو	Explain in detail about Normal Forms	(03
8	With Suitable examples.	

Mark Allocation

Rubrics	Marks Allocated	Marks obtained
Content Quality	6	4
Presentation Quality	2	2
Timely submission	2	1
Total marks	10	7

Name and Signature of the Faculty Incharge

ToD/CSE

Assignment - 2 Data Base Management System

Name: B. Saran

Rept : BECSE

Sub Code: CS 8492

Regno: 73242210/4038

Jan Jan

Moronal Johns ! * Normalization is the process of orangonizing relate in ne date base so that it means two lasic vaguerements! 1) There is no reductancy of whata 2) Data depondencies rare logical The normalization is important because it villaco detalase to take less space. * It roles helps in increasing the performance first Normal form: Fre table is said to be in INF if it follows i) It should only have single Catomic) valued attributes 2) Values stored in a table should be of the Same domain 3) All-the volumn should be have unique names a) And the order in which data is stored, does not matter

Snine	Phone
AAA	(1111)
BBB	22220
ccc	8333
2.5	7444
	BBB

As there were multiple values of phono number for Sid I and 3 , the salvove table is not in INE We make it in INF. The conversion is a follow

Sid	Somme	Phone
	AAA	11 11
a	BBB	2222
3	ccc	3933
<u>چ</u>	CCC	4444
4	DDD	5555

Second Normal John:

Before understanding the second normal form to us first discuss the concept of partial functional dependency and prime and non prime relations Concept of Partial functional Dependency Partial dependency means that a non prime

sattailanter is functionally adependent on par of in candidate tog. Holgenouple: Consider on Adulian RCNB, C,O) with Sunctionally Dependent [AB> CD>) N>C] Here (AB) is un sandidate Loy Cocauso (A) = [ABCN]= PR] Burne and Non Prime Attributer * Psino Attilato: An attilate, which is a part of the cardidate key, is known is a paine attribute Non Brime Altribute: An altribute, which is not so part of the primary key, is said to be non paine altibute Example: Consider a Relation R= JA, B, C, D} Candidate They as AB The Second Normal form! For a table to be in the Second normal form, following constition must be followed. i) It should be in first normal form II should not have partial functional dependency Jos example: Consider a table with every

Student Course

Sid	Smme	ved	Chamo
,	AAA	101	C
2	BBB	102	C-+++
3	CCC	loi	C
4	DDD	103	Java

JRs table is not in 2NF. For converting rations table to DNF we must follow the following the Step 1: The valore table is in . INF Step 2! Here & name and sid were rassociated Similarly cid and Crame are ussociated with each other. Now if we delete a record with sides then Automatically the course cost will ralso get • rdeleted: Student

4 ·1	/ /	1
Sid/	Snamo	Cid
Y	AAA	tol
2	BBB	102
3	ccc	(02
21	DDD	1021

philse

Cid	· Chame
101.	C
102	C++
loi	C
102	Jara

Third Normal Form:

import of Super key and vandidate key import key! A Super is a set or some of more. Clumns. to uniquely identify, hows in a table indialate key! The minimal set of attailante which can uniquely indentify a tuple is known as

Candidate Bay.

RegID	Rollno	Sname
101		AAA
102	a	BBB
103	a	ccc
104	7	DDD

Superbays

* [RegID]

* [Reg ID, Roll no]

* [Reg ID , Sname]

* (Roll No, Grame) * LRegID, Rollno, Sname Candidato Kay! * [Reg Io] * [Roll no] Third Normal Form! A table is said to be in the Shird Normal form When, 1) It is in the Second Normal form 11) It does not have transtive Lapandency For example: Consider following table Student_ details ous follows. Sid Sname Zipcole Cityrame AAA 1411 Pune BBB 2222 Surat 3333 Ccc Jamilnadu 4444 Rajustan DDDD 21 Superkey! [8id], [Sid Sname], [Sid Sname, ziplote [3id, Zip Code, atyname]. and Son, Candidate Rey: [Sid]

Student

	the second control of		
Com one of Lapshing	Sid	Sname	Zipado
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Capital and Stages	Q	BBB	2222
GD()(disc)(2)(disc)(disc)	8.	Ccc	3223
A Company of Street, S	1	DDD	2444

Zip

Zip Cole	City have
1.11	Rune
2222	Swot
8333	Channai
4444	Joi pur



Agenda

- · What is Data?
- Types of Data
- Structured Data and Unstructured Data
- How do we store Data into Computers ?
- File System Vs Database System
- File System Issues
- What is Management?
- What is DBMS?

What is Data?

Factual information (such as measurements or statistics) used as a basis for reasoning, discussion, or calculation

From Wikipedia:

Data are characteristics, usually numerical, that are collected through observation. In a more technical sense, data is a set of values of qualitative or quantitative variables about one or more persons or objects, while a datum (singular or data) is a single value of a single variable.

• Example: Name, Address, Zip, SSN are characteristics or information about a person.

Types of Data

- Unstructured-Data
- Structured Information (interpreted data data supplied with semantics)

Structured and Unstructured Data

Structures data is most often categorized as quantitation data, and its the type of data most of via are used to exemply with. Extra of data that fits mestly within Excellence and powers in relational databases and spressionness.

Examples of structured data include narries, dates, additional practicular numbers, stock information, protocolation, and more.

Uniformitation distalls moved offers callegoe over all our sofetime state, and it commits be proported and sharpferd out by surpaintitional losses and inventors.

How do we store Data into Computers?.

- Data could be stored computers in File System and or Database Management and or content management systems.
- * Let us focus on File System and Database approaches , Each one has its own advantages and disadvantages.

File System Vs Database System

File System

Database System



File System Vs Database System Cont...

File System

Database System



File System Issues

- Redundancy of data: Data is said to be redundant if same data is copied at many places, if a student wants to change Phone number, he has to get it updated at various sections. Similarly, old records must be deleted from all sections representing that student.
 Inconsistency of Data: Data is add to be inconsistent if multiple copies of same data does not match with each other if Phone number is different in Accounts Section and Academics Section, it will be inconsistency. Inconsistency may be because of typing errors or not updating all copies of same data.
 Difficult Data Access: A user should know the exact location of file to access data, so the process is very cumbersome and tedius. If user wants to rearch student hostel allotment number of a student from 10000 insorted students' records, how difficult it can be considered access. If it is system may lead to unauthorized access to data if a student gets access to file having his marks, he can change if in unauthorized access to data if a student gets access to file having his marks, he can change it in unauthorized access.

- No Concurrent Access: The access of same data by multiple users at same time is known as concurrency. File system does not allow concurrency as data can be accessed by only one user at a time.
 No Backup and Recovery: File system does not incorporate any backup and recovery of data if a file is list or corrupted.

What is Management?

- Generally Management refers create, Retrieve, update and delete.
- For example: Money Management refers earning the money, querying money, distribute the money and spend the money.

What is Database?

- Database: Database is a collection of inter-related data which helps in
 efficient retrieval, insertion and deletion of data from database and
 organizes the data in the form of tables, views, schemas, reports etc. For
 Example, university database organizes the data about students, faculty,
 and admin staff etc. which helps in efficient retrieval, insertion and
 deletion of data from it.
- Now Database Management talks about
 Create Database
 Retrieve(View) database
 Update (database
 Delete Database

What is DBMS?

- Database Management System: The software which is used to manage database is called Database Management System (DBMS). For Example, MySQL, Oracle etc. are popular commercial DBMS used in different applications. DBMS allows users the following tasks:
- Data Definition: It helps in creation, modification and removal of definitions that define the organization of data in database.
- Data Updation: It helps in insertion, modification and deletion of the actual data in the database.
- Data Retrieval: It helps in retrieval of data from the database which can be used by applications for various purposes.
- User Administration: It helps in registering and monitoring users, enforcing data security, monitoring performance, maintaining data integrity, dealing with concurrency control and recovering information corrupted by unexpected failure.

Database Languages

· A database system provides a data definition language to specify the database schema and a data manipulation language to express database queries and updates. In practice, the data definition and data manipulation languages are not two separate languages; instead they simply form parts of a single database language, such as the widely used SQL language

Data Definition Language(DDL)

We specify a distribute schedule by a ser of definitions expressed by a special language whiled a deas definition language (DDI). For instance, the following elaboration in the NGS - language defines the environmental language and the expression of the definition of the definition of the definition of the above 100. As the entire contact the environmental language and environmental language and the definition of the definition of explanation of the above 100. As the entire to entire the first of the definition of explanation of the environmental language and environmental language

modelying actival data.
We speech the showard enteriors and secure methods could be the database explain by a set at interactive is a special type of 1501, whiled a data staturage and definition from the statements in a special type of 1501, whiled a data statement and definition from the statements of the statements of the statements of the statements of the database schements which are smaller badden booker to some it. The data values, where it is a statement of the database point stated, certain consistency constitutions for examples suppose the battone on an incommon should not belt below 3700. The DML gravation of the database is appeared to the database is appeared.

Data Manipulation Language(DML)

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The critical information insend in the distribute.

The critical in information is not the distribute.

The defence of information provides the starbute.

The defence of information provides the starbute.

The defence of information information that distribute.

A data manipolation language (DML) is a language that modifies some as access on moripidate date as regardered by the appropriate dark model. There are baseably from

Procedural DMLs require a user to specify solest data are montrel and have to get those

Acts

Destination DMLs calso referred to a sample content at DMLs is express a since to specify only the data are neighbor specifying from the art of since to superate or both data are neighbor specifying from the art most of extra procedural DMLs. Destination DMLs are mostly ensure to learn and use that one procedural DMLs. However, since a most does not have to appearly from to get the data, the database system has to figure one one officially an ensure of according that. The DML companion of the NQL language is neceporately.

DML Continuation...

A query is a statement sequenting the retrieval of information. The portion of a DML dist months information retrieval is called a query language. Although technically incorrect, if in insumon practice to use the testin query language and data montpulation kinguage.

system) memby.

This query in the 5QL language finels for name of the customer where customered.

ried emission existence many

where customer concerner of = 192-83-7465

The query spectus, that those areas from the table entremes where the entremested to 1976-0-769 must be entremed, and the contract areas attached of those twice must be

symmetric may make be arranged, and the company many attribute of those tries a mast be displayed.

General stay break a information from mose than one table. For metasive the fulls a sing-group facels, the ballouse of all decounts around by the informer with a intercept of 122-s.t.

DML Continuation...

welect account halance

В ни фрики, отому

where depositor customer of = 192-83-7465 and

деромы осеран натые « ассилы ассилы азавы

There are a number of disabase query lineranges to use, asther commercially or

The levels of abstraction apply not only to defining of structuring data, but also to manapoliting data. At the physical level, we must define algorithms that allow efficient access to data. At higher levels of abstraction, we emphasize case of use. The goal is to allow laminum to interact efficiently with the system. The query processed compensation the database system remisers DML queries into sequences of setions at the physical level of the database system.

Database Management Systems UNAT-I Short answer Questions

- 1. Define the terms data and information?
- 2. Define (i) Database (ii) DBMS
- 3. List the advantages and applications of DBMS?
- 4. What are the disadvantages of file processing system?
- 5. Define instances and schemas of database?
- 6. What is data model? List the types of data models?
- 7. Discuss about Data Definition language?
- 8. Discuss about Data Manipulation language?
- 9. What is data Abstraction? Give the levels of data abstraction?
- 10. Who is DBA? What are the responsibilities of DBA?
- 11. Discuss Data Independence?
- 12. What is an entity relationship model?
- 13. Define (i) Entity (ii) Attribute
- 14. Define Relationship and Relationship set?
- 15. What are key constraint and participating constraints?
- 16. Define weak entity and strong entity sets?
- 17. Define relation, relation instance and relation schema.
- 18. Define i) super key ii)candidate key iii) primary key
- 19. Explain the use of foreign key constraint?
- 20. Define the terms arity and cardinality of relation?
- 21. What are domain constraints
- 22. Explain about querying relational data?
- 23. Define views?
- 24. Discuss how can you change the data in the table?
- 25. List various types of attributes?
- 26. Discuss how can you alter and destroy tables?
- 27. Explain the use of null values?

- 1. Compare and Contrast file Systems with database system?
- 2. Define Data Abstraction and discuss levels of Abstraction?
- 3. Discuss about different types of Data models?
- 4. Describe the architecture of DBMS?
- 5. Discuss additional features of the ER-Models?
- 6. Discuss about the Conceptual Design with the ER-Model?
- 7. Write about views and updates on views?
- 8. Explain different types of database users and write the functions of DBA?
- 9. Explain about different types of integrity constraints?
- 10. Discuss about the logical database Design?
- 11. Distinguish strong entity set with weak entity set? Draw an ER diagram to illustrate Weak entity set?
- 12. Explain how the integrity constraints are specified and enforces?
- 13. Explain in detail about views?

UNIT-II Short answer Questions

- 1 Define relational database query?
- 2 Explain different types of query languages?
- 3 Explain about relational algebra?
- 4 State about SELECT operation in Relational algebra?
- 5 State about PROJECT operation in Relational algebra?
- 6 Explain about set operations?
- 7 Discuss the use of rename operation?
- 8 Define join? Explain different join operations?
- 9 Illustrate division operation?
- 10 Explain about tuple relational calculus?
- 11 Explain about Domain relational calculus?
- 12 Discuss about the expressive power of relational algebra and calculus?

- 13 Discuss the basic form of SOL query?
- 14 Explain the working of union, intersection and except operations?
- 15 Define nested queries?
- 16 Define correlated nested queries?
- 17 Explain Aggregate Functions?
- 18 What is the use of groupby and having clauses?
- 19 Define Null Values?
- 20 Define tuple variable with its syntax?
- 21 Define outer join? Explain its types?
- 22 Explain how to create new domain?
- 23 Define Assertions?
- 24 Discuss about trigger?
- 25 Demonstrate how to add a NOT NULL column to a table?
- 26 Write a TRC query to find the names of sailors who have reserved boat 103?
- 27 Write a DRC query to find the names of sailors who have reserved red boat?

- 1. Illustrate different operations in Relational algebra with an example?
- 2. Define Join? Explain different types of joins?
- 3. Discuss about Relational calculus in detail?
- 4. Define trigger and explain its three parts? Differentiate row level and statement level triggers?
- 5. Illustrate Group by and having clauses with examples?
- 6. Discuss about Complex integrity constraints in SQL?
- 7. Define null value? Describe the effect of null values in database?
- 8. Discuss different types of aggregate operators with examples in SQL?
- a Define a nested query?
 - b Write a nested query to find the names of sailors who have reserved both a red and green boat?
 - c. Write a nested query to find the names of sailors who have reserved all boats?

- 9 a. Discuss correlated nested queries?
 - b. Write a query to find the names of sailors who have reserved a red boat?
- c. Write a query to find the names of sailors who have not reserved a red boat?

UNIT-III

Short answer Questions

- Define redundancy?
- 2. Define functional dependency?
- Explain the problems with Redundancy? 3.
- What is decomposition? Explain the properties of Decomposition? 4.
- Discuss normalization? 5.
- Illustrate functional dependency with example?
- Illustrate fully functional dependency with example?
- Demonstrate transitive dependency? Give an example? 9.
- Define First Normal Form?
- 10. Define Second Normal Form?
- 11. Define Third Normal Form?
- 12. Explain about Loss Less Join Decomposition?
- 13. Describe Dependency Preserving Decomposition?
- 14. What is multi valued Dependency?
- 15. Define Fourth Normal Form?
- 16. Define Join Dependency?
- 17. Define BCNF?
- 18. Explain Fifth Normal Form?
- 19. Explain about Inclusion Dependency?

Long answer Questions

1. Illustrate redundancy and the problems that it can cause

- 2. Define decomposition and how does it address redundancy? Discuss the problems that may be caused by the use of decompositions?
- 3. Define functional dependencies. How are primary keys related to FD's?
- 4. Define normalization? Explain 1NF,2NF,3NF normal forms
- 5. Compare and contrast BCNF with 3NF?
- 6. Describe properties of decompositions

UNIT-IV

Short answer Questions

- 1. Define a Transaction? List the properties of transaction
- 2. Discuss different phases(states) of transaction?
- 3. What is shadow copy technique?
- 4. List the advantages of concurrent execution?
- 5. Define Schedule? What is a serial schedule?
- 6. Discuss the Procedure to test Serializability?
- 7. Demonstrate Conflict Serializability?
- 8. Discuss View Serializability?
- 9. Discuss recoverable schedules?
- 10. Discuss cascade less schedules?
- 11. Explain the procedure to test for serializability?
- 12. Explain about different types of locks?
- 13. Define Deadlock?
- 14. Explain about locking protocols?
- 15. Define Two Phase locking protocol?
- 16. Demonstrate the implementation of Isolation?
- 17. Explain how the locks are implemented?
- 18. Explain the rules of tree protocol?
- 19. What is timestamp? Explain different timestamps used by a transaction?
- 20. Explain Thomas write rule?
- 21. What are the phases of validation based protocol?
- 22. Explain different timestamps used by validation protocol?
- 23. Define granularity?

- 24. Discuss about Failure Classification?
- 25. What are different storage types?
- 26. What are the fields of update record?
- 27. Discuss log based recovery?
- 28. Differentiate deferred and immediate database modifications?
- 29. Define a checkpoint?
- 30. Discuss the failures that can occur with loss of Non-volatile storage
- 31. Explain about ARIES?

- 1. Explain ACID properties and illustrate them through examples?
- 2. Discuss How do you implement Atomicity and Durability
- 3. Illustrate Concurrent execution of transaction with examples
- 4. Discuss Serializability in detail?
- 5. Discuss two phase locking protocol and strict two phase locking protocols?
- 6. Describe Times tamp based locking protocols?
- 7. Describe Validation-based locking protocols?
- 8. Discuss in detail Multiple Granularity?
- Explain in detail storage structure
- 10. Discuss how do you recover from failure?
- 11. Explain Buffer Management?
- 12. Explain different types of advanced recovery techniques
- 13. Write in detail about Remote Backup systems?

UNIT-V

Short answer Questions

- 1. Discuss about data on External storage?
- 2. What is indexing and what are the different kinds of indexing?
- 3. Explain Clustered Indexes?
- 4. Discuss the Primary and Secondary indexes?
- Define Tree Indexing?
- 6. Explain Hash based Indexing?
- 7. Compare different file organizations?
- 8. Discuss the intuition for Tree Indexes?
- 9. Define Indexed Sequential Access Method?
- 10. Discuss about Overflow pages and Locking considerations of ISAM?
- 11. Discuss the Cost model of Heap files, Sorted files and Clustered files?
- 12. Explain the structure of B+ tree?
- 13. Describe how the insert and delete operations are performed in B+ tree?
- 14. Explain how search is performed in B+ tree?
- 15. Define static Hashing?
- 16. Explain extendible hashing?
- 17. Define linear hashing?
- 18. Differentiate between linear and extensible hashing?

- 1. Write in detail about hash based indexing and Tree based indexing
- 2. Compare I/O costs for all file organizations
- 3. Explain in detail about ISAM
- 4. Explain about B+ tree index file?
- 5. Demonstrate searching a given element in B+ trees? Explain with example?
- 6. Illustrate insertion of an element in B+ Tree with example
- 7. Illustrate deletion of an element in B+ Tree with example
- 8. Write in detail about Static Hashing
- 9. Explain in detail about Extendible hashing
- 10. Explain in detail about Linear hashing

11. Compare and contrast Extendible hashing With Linear hashing

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Reg. No.

Question Paper Code: 80099

Tech. DEGREE EXAMINATIONS, APRIL/MAY 2019.

Third/Fourth Semester

Information Technology

CS 8492 + DATABASE MANAGEMENT SYSTEMS

(Cominon to Computer Science and Engineering/Computer and Communication Engineering)

(Regulation 2017)

Time: Three hours

Maximum: 100 marks

Answer ALL questions.

PART A — $(10 \times 2 = 20 \text{ marks})$

- 1. What is a data model? List the types of data model used.
- 2. List any eight applications of DBMS.
- 3. Give the properties of decomposition.
- 4. Define the terms Entity set and Relationship set.
- 5. What are the states of transaction?
- 6. What is meant by log-based recovery?
- 7. Define dense index.
- 8. Mention all the operations of files.
- 9. Mention two features of Multimedia databases
- 10. Compare sequential access devices versus random access devices with an example.

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PART B + (5 x 13 = 65 marks)

11. (a) Explain the three different groups of data models with suitable axemples

Or

- (b) Describe about the static and dynamic SQL in detail.
- 12 (a) What is normulization? Explain in detail about all Normal forms.

Or

- (b) Briefly discuss about the functional dependency concepts.
- 13. (a) Discuss in detail about the testing of serializability.

Or

- (b) Explain deferred and immediate modification versions of the log based recovery scheme.
- 14. (a) What is RAID? Briefly discuss about RAID

Or

- (b) Describe the structure of B+ tree and give the algorithm for search in the B+ tree with example.
- 15 (a) Discuss in detail about the distributed databases.

Or

(b) Explain in detail about the Deductive DB and Spatial DB.

PART $G \leftarrow (1 \times 15 = 15 \text{ marks})$

(Application / Design / Analysis / Evaluation / Creativity/Case Study questions)

16 (a) Discuss in detail about the ACID properties of a transaction

or or the

(b) What is concurrency control? How it is implemented in DBMS? Briefly elaborate with suitable diagrams and examples.

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x) Departments have a professor (known as the chairman) who runs the department

- Professors work in one or more departments, and for each department that they work in, a time percentage is associated with their job-
- xii) Graduate students have one major department in which they are working
- in in les a student advisor) who advisos him or her on what courses to take xiii) Each graduate student has another, more senior graduate student known

Design and draw an ER diagram that captures the information about the university.

Be sure to indicate any key and participation constraints. Use only the basic EE model here; that is, entities, relationships and attributes. (5+10)

For the following relation schema R and set of functional dependencies F: R(A, B, C, D, E), $F = \{AC \to E, B \to D, E \to A\}$. List all candidate keys. 69

E Consider the Table-16 and answer to queries given below.

Table-16 User_personal.

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- i) Is this table in First Normal Form-INF? Justify and normalize to INF if
- 2) Je this table in Second Normal Form-2NF 2 Justify and normalize to 2NF
- بن Is User personal in Third Normal Form-3NF? Justify and normalize to If needed

Question Paper Code: 90156

Tech DEGREE EXAMINATIONS, NOVEMBER/DECEMBER 2019 Third/Fourth/Fifth Semester

📆 8492 – DATABASE MANAGEMENT SYSTEMS Computer Science and Engineering

(Common to Computer and Communication Engineering /Mechanical and Automation Engineering/Information Technology)

(Regulations 2017)

Maximum: 100 Marks

Time: Three Hours

Answer ALL questions

PART-A

(10×2=20 Marks)

l. What are the four main characteristics that differentiate the database approach from the file-processing approach?

- 2. Express in relational algebra, the division operation(/) using the project, cartesian product and minus operations. Give a simple example.
- .3. . Boyce-Codd normal form is found to be stricter than third normal form'. Justify
- -4: What is the significance of "participation role name" in the description of relationship types?
- ... 5...List the responsibilities of a DBMS has whenever a transaction is submitted to the system for execution?
- 6. Brief any two violations that may occur if a transaction executes a lower isolation
- How do you represent leaf node of a B' tree of order p?
- Which cost components contribute to query execution?
- List information types of documents necessary for relevance ranking of documents
- 10. What one could understand from allocation schema?

process of relational database design. (7) process of relational database design. (8) process of relational database design. (OR) (OR) b) 1) Explain with suitable example, the constraints of specialization and generalization in ER data modeling. (7) generalization in ER data modeling. (6)	(OR) b) B Consider the schema given in question no. 11.a) i) and write the following queries in SQL 1) Find the names of suppliers who supply some red part. 2) Find the sids of suppliers who supply some red part. 3) Find the sids of suppliers who supply every red part. 4) Find the pids of parts supplied by atleast two different suppliers. (5) Explain the three schema architecture with a neat diagram.	Write the following queries in relational algebras. 1) Find the sids of suppliers who supply some red or green part. 2) Find the sids of suppliers who supply some red part or are at 221 Packer. Street. 3) Find the pids of parts supplied by atleast two different suppliers. ii) Sketch the typical component modules of DBMS. Indicate and explain the interactions between those modules of the system. (7)	Suppliers (sid: integer, sname: string, address: string) Parts (pid: integer, pname: string, color: string) Catalog (sid: integer, pid: integer, cost: real) The key fields are underlined and the domain of each field is listed after the field name. Therefore sid is the key for Suppliers, pid is the key for Parts and sid and pid together form the key for Catalog. The Catalog relation lists the parces charged for parts by suppliers. (6)	90156 PART - B ACCESS 3.200 Study Materials for Sen (5x13=65 Marks)
co-investigators). vi) Professors can manage and/or work on multiple projects. vii) Professors can manage and/or work on multiple projects. vii) Each project is worked on by one or more graduate students (known as the project's research assistants). When graduate students work on a project, a professor must supervise their work on the project. Graduate students can work on multiple projects, in work on the project. Graduate students can work on multiple projects, in which case they will have a (potentially different) supervisor for each one. ix) Departments have a department number, a department name and a main office.	PART_C. (1×15=15 Marks) 16. a) Consider the following information about a university database: 1) Professors have an SSN, a name, an age, a rank and a research specialty. 1) Projects have a project number, a sponsor name (e.g., NSF), a starting date, an ending date and a budget. an ending date and a budget. iii) Graduate students have an SSN, a name, an age and a degree program (e.g., M.S. or Ph.D.). (e.g., M.S. or Ph.D.). iv) Each project is managed by one professor (known as the project's principal investigator).	b) i) With simple algorithms explain the companies: Block Nested-loop join. (3) Block Nested-loop join. (3) Block Nested-loop join. (3) Sketch and concise the basic steps in Query Processing. (9) 15. a) i) Illustrate the usage of OQL, the DMG's query language. (4) ii) Brief on the methods to store XML documents. (OR) (OR) (OR) (9) b) i) Illustrate the approaches to store relations in distributed database. (9) ii) How effectiveness of retrieval is measured? Discuss. (4)		13. a) Discuss claborately the two-phase locking protocol that ensures serializability. (4)